

Exploring the abyss in Reverse Mathematics

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Secondly, we discuss the **very different** **logical** and **mathematical** limits of this extension. (chasm, abyss)

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Secondly, we discuss the very different logical and mathematical limits of this extension. (chasm, abyss)

Thirdly, we may discuss foundational implications, though ...

Reverse Mathematics

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= finding the **minimal** axioms needed to **prove** a theorem of ordinary mathematics.

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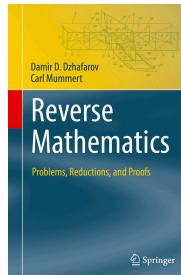
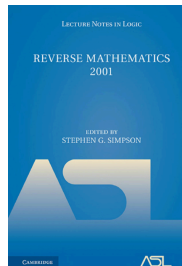
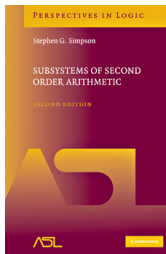
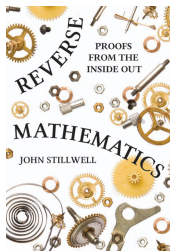
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Harvey Friedman & Steve Simpson (courtesy of MFO).

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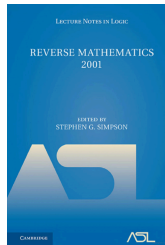
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Mummert: a few equivalences for Π_2^1 -comprehension and topology.

Higher-order Reverse Mathematics

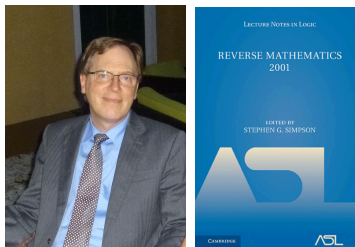
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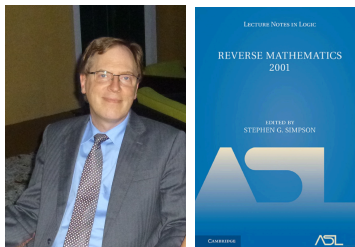
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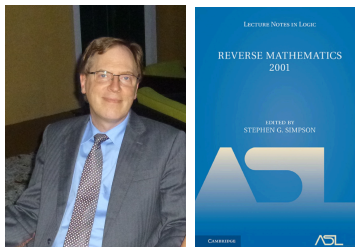
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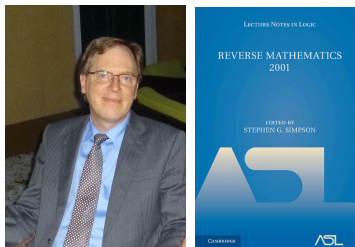


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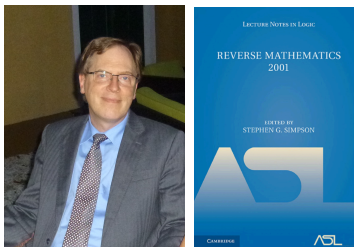


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Kohlenbach's higher-order RM uses the richer **language of all finite types**. Thus, the **use of codes or representations is seriously reduced**. E.g. discontinuous functions on \mathbb{R} are directly available.

The base theory RCA_0^ω

RCA_0^ω makes use of the **language of finite types**: $n \in \mathbb{N}$ or n^0 ,
 $f \in \mathbb{N}^{\mathbb{N}}$ or f^1 , $Y : \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}$ or Y^2 , et cetera.

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Real numbers and ' $=_{\mathbb{R}}$ ' defined **as in RCA_0** ; $\mathbb{R} \rightarrow \mathbb{R}$ -functions are $\mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}^{\mathbb{N}}$ -functions extensional relative to ' $=_{\mathbb{R}}$ '.

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- **third-order** theorems about (slightly) discontinuous functions.

These **third-order theorems** are called **second-order-ish** for obvious reasons. A similar phenomenon does **not** exist for first- and second-order theorems (AFAIK).

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- Cousin's lemma for **continuous** functions.

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The following **third-order thms** are equivalent to WKL_0 over RCA_0^ω :

- A **regulated** function on $[0, 1]$ is bounded.
- A bounded **Baire 1** function on $[0, 1]$ has a supremum.
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WILD: there are 2^c **non-measurable** quasi-continuous functions and 2^c **non-Borel** bounded and measurable quasi-continuous functions.

Arithmetical comprehension

The following are equivalent to ACA_0 over RCA_0 :

- Let $F : C \rightarrow \mathbb{R}$ be **continuous** where $C \subset [0, 1]$ is an RM-closed set. Then $\sup_{x \in C} F(x)$ exists.
- Let $F : C \rightarrow \mathbb{R}$ be **continuous** where $C \subset [0, 1]$ is an RM-closed set. Then F attains a maximum value on C .
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These **third-order thms** are equivalent to ACA_0 over RCA_0^ω :

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These **third-order thms** are equivalent to ATR_0 over RCA_0^ω :

- Jordan decomposition theorem restricted to **arithmetical** (or: Σ_1^1) functions.
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Baire (1905) notes that **Baire 2 functions can be represented as iterated limits**.

Π_1^1 -comprehension

These **third-order thms** are equivalent to Π_1^1 -CA₀ over $RCA_0^\omega + X$:

- For any $x \in \mathbb{N}^{\mathbb{N}}$, any bounded $\Sigma_1^{1,x}$ -class in \mathbb{Q}^+ has a supremum.
- A bounded **effectively Baire 2** $f : [0, 1] \rightarrow \mathbb{R}$ has a supremum.
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Effectively Baire 2 means: iterated limit of **double** sequence of continuous functions (\approx second-order codes for Baire 2).

Baire (1905) notes that **Baire 2 functions can be represented as iterated limits**.

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There are however **hard limits** to the Biggest Five phenomenon, with interesting consequences.

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The above was obtained based on the RM of Kleene's \exists^2 :

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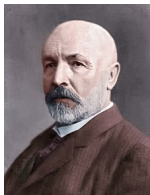
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The following is not provable in $\text{RCA}_0^\omega + (\exists^2) + \text{Z}_2$:

There is a $\mathbb{R} \rightarrow \mathbb{R}$ -function that is **not Baire 2**.

Exploring the abyss: the uncountability of \mathbb{R}

Cantor's first set theory paper (1874): **uncountability of \mathbb{R}** .



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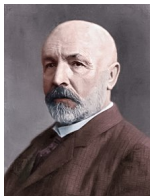


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Many many many (third-order) mainstream theorems imply NIN or NBI. **However**, NIN and NBI **cannot** be proved in $\text{RCA}_0^\omega + \text{Z}_2$ and stronger (higher-order) systems (see Normann-Sanders, JSL, 2022).

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Many equivalences for NIN and basic properties of **regulated** functions.

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The theorems in (b) are called **second-order-ish** for obvious reasons.

There are a **gazillion** possible equivalences, warranting the name **the Biggest Five phenomenon**.

Slight variations or generalisations of the theorems in (b) imply NIN and cannot be proved in $\text{RCA}_0^\omega + \mathbb{Z}_2$ and stronger systems.

Similar results for **WWKL**, Vitali's covering lemma, and Kleene's (\exists^2).

Many equivalences for NIN and basic properties of **regulated** functions. Same for basic properties of measure and category and **semi-continuity** (Baire, Volterra, ...).

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Thanks!
Questions?

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