# Weak König's Lemma in the absence of $\Sigma^0_1$ induction

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## Background: reverse mathematics

Reverse mathematics studies the strength of axioms needed to prove mathematical theorems. This is done by deriving implications between the theorems and set existence principles expressed in the language of second-order arithmetic, which has:

- $\blacktriangleright$  vbles x, y, z, . . . , i, j, k . . . for natural numbers,
- vbles X, Y, Z, ... for sets of naturals,
- ▶ non-logical symbols  $+, \cdot, 2^x, \leq, 0, 1, \in$ .

Often, the theorems studied are  $\Pi_2^1$  of the form  $\forall X \exists Y \psi$ , and their strength is related to the difficulty of computing Y given X.

The implications are proved in a relatively weak base theory.

## The usual base theory, and an important axiom

The usual base theory, RCA<sub>0</sub>, has the following axioms:

- $\blacktriangleright$  +,  $\cdot$ , 2<sup>x</sup> etc. have their usual basic properties.
- $\Delta_1^0$  comprehension: if  $\bar{X} = X_1, \dots, X_k$  are sets and  $\psi(x, \bar{X})$  is computable relative to  $\bar{X}$ , then  $\{n : \psi(n, \bar{X})\}$  is a set.
- ▶  $\Sigma_1^0$  induction: if  $\overline{X}$  are sets and  $\psi(\mathbf{x}, \overline{X})$  is c.e. relative to  $\overline{X}$ , then  $\psi(\mathbf{0}, \overline{X}) \land \forall n (\psi(n, \overline{X}) \Rightarrow \psi(n + 1, \overline{X})) \Rightarrow \forall n \psi(n, \overline{X})$ .

Possibly the most important theory in reverse mathematics, WKL<sub>0</sub>, is axiomatized by RCA<sub>0</sub> and Weak König's Lemma WKL: "Every infinite tree in  $\{0, 1\}^{\mathbb{N}}$  has an infinite path".

This says essentially: "The interval [0, 1] is Heine-Borel compact". Or "For every set X there is a set Y of PA degree relative to X".

## Properties of WKL

 $RCA_0$  proves:  $WKL_0 \equiv a$  plethora of theorems, from compactness of first-order logic to the Peano existence thm for ODE's.

WKL is not provable in  $RCA_0$ . On the other hand:

Theorem (Harrington 1977, independently Ratajczyk 1980's) WKL<sub>0</sub> is  $\Pi_1^1$ -conservative over RCA<sub>0</sub>, i.e. every  $\Pi_1^1$  sentence provable in WKL<sub>0</sub> is also provable in RCA<sub>0</sub>.

The proof is by adding a path to an infinite 0-1 tree T in a countable model of  $RCA_0$ , which is done via forcing with infinite subtrees of T.

## A weaker base theory

In an alternative, weaker base theory RCA<sub>0</sub><sup>\*</sup>, one replaces  $\Sigma_1^0$  induction with  $\Delta_1^0$  induction: if  $\bar{X}$  are sets and  $\psi(x, \bar{X})$  is computable relative to  $\bar{X}$ , then  $\psi(0, \bar{X}) \land \forall n (\psi(n, \bar{X}) \Rightarrow \psi(n + 1, \bar{X})) \Rightarrow \forall n \psi(n, \bar{X}).$ 

- Used to identify some theorems that are equivalent to Σ<sub>1</sub><sup>0</sup> induction (e.g. "every non-zero poly has finitely many roots" [Simpson-Smith 1986]) and some that do not need it (e.g. Friedberg-Muchnik Thm [Chong-Mourad 1992]).
- Turns out to be useful in understanding some aspects of reverse mathematics over RCA<sub>0</sub> (e.g. [Belanger 20XX]).

## Weak König's Lemma over RCA<sub>0</sub>\*

- The theory obtained by adding WKL to  $RCA_0^*$  is known as  $WKL_0^*$ .
- Theorem (Simpson-Smith 1986) WKL<sub>0</sub><sup>\*</sup> is  $\Pi_1^1$ -conservative over RCA<sub>0</sub><sup>\*</sup>.

The proof is a forcing argument similar to the one over RCA<sub>0</sub>.

However: today's talk is about a property that models of  $\mathsf{WKL}_0^*+\neg \mathsf{I}\Sigma_1^0$  have but those of  $\mathsf{WKL}_0$  do not.

## The isomorphism theorem

#### Theorem

Let  $(M, \mathcal{X})$ ,  $(M, \mathcal{Y})$  be countable models of WKL<sub>0</sub><sup>\*</sup>, and assume that  $(M, \mathcal{X} \cap \mathcal{Y}) \models \neg I\Sigma_1^0$ . Then  $(M, \mathcal{X}) \simeq (M, \mathcal{Y})$ .

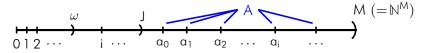
- This can be seen as a second-order generalization of results due to Kossak and Kaye saying that models of IΔ<sub>1</sub> + ¬IΣ<sub>1</sub> have "many" automorphisms.
- There are many ω-models of WKL<sub>0</sub> ("Scott sets") that are neither isomorphic nor elementarily equivalent to one another.
- Any countable  $(M, W) \models \text{RCA}_0$  will have extensions  $(M, \mathcal{X}), (M, \mathcal{Y})$  satisfying WKL<sub>0</sub> with  $(M, \mathcal{X}) \not\equiv (M, \mathcal{Y})$ .

## Plan for rest of talk

- Brief comment about the proof of the isomorphism theorem
- Consequences for WKL<sub>0</sub><sup>\*</sup> in the absence of  $I\Sigma_1^0$ .
- Consequences for reverse mathematics over RCA<sub>0</sub>.

### The isomorphism theorem: ideas behind proof

Because (M, X ∩ Y) satisfies ¬lΣ<sub>1</sub><sup>0</sup> there is a Σ<sub>1</sub><sup>0</sup>-definable proper cut J closed under x → 2<sup>x</sup> and an infinite set A ∈ X ∩ Y s.t. A = {a<sub>i</sub> : i ∈ J} enumerated in increasing order.



We use back-and-forth. At each step, we have finite tuples r̄, R in the domain, s̄, S̄ in the range of the partial iso. The invariant is roughly: for each Δ<sub>0</sub> formula δ, each i, k ∈ J,

$$(\mathsf{M},\mathcal{X})\models\delta(\mathsf{a}_i,k,\bar{r},\bar{R}) \text{ iff } (\mathsf{M},\mathcal{Y})\models\delta(\mathsf{a}_i,k,\bar{s},\bar{S}).$$

WKL needed to find "globally good" second-order element to add in the inductive step, based on "locally good" ones that are easier to find directly from inductive assumption.

## Consequences for WKL<sub>0</sub><sup>\*</sup>: the analytic hierarchy

For a set W, let  $W_k = \{n : \langle k, n \rangle \in W\}.$ 

If  $(M, \mathcal{X}) \models WKL_0^*$  and  $A \in \mathcal{X}$ , then there exists  $W \in \mathcal{X}$  such that  $W_0 = A$  and  $(M, \{W_k : k \in M\}) \models WKL_0^*$ . We say that W codes a model of  $WKL_0^*$ . If it satisfies  $\neg I\Sigma_1^0$ , then by the isomorphism theorem it is elementarily equivalent to  $(M, \mathcal{X})!$ 

#### Corollary

For any formula  $\psi(\bar{x}, \bar{X})$ , TFAE provably in WKL<sub>0</sub><sup>\*</sup> +  $\neg I\Sigma_1^0$ :

- (i)  $\psi(\bar{\mathbf{x}},\bar{\mathbf{X}})$ ,
- (ii) "there exists a coded model of WKL\_0^\* +  $\neg I\Sigma_1^0 + \psi(\bar{x},\bar{X})$ ",
- (iii) "there is no coded model of WKL\_0^\* +  $\neg I\Sigma_1^0$  +  $\neg \psi(\bar{x},\bar{X})$ ".

Thus, in WKL\_0^\* +  $\neg l \Sigma_1^0$  the analytic hierarchy collapses to  $\Delta_1^1.$ 

## Consequences for WKL<sub>0</sub><sup>\*</sup>: conservativity

#### Corollary

Let  $\psi$  be a  $\Pi_2^1$  statement. Then:

(i) ψ is Π<sub>1</sub><sup>1</sup>-conservative over RCA<sub>0</sub><sup>\*</sup> + ¬IΣ<sub>1</sub><sup>0</sup> iff WKL<sub>0</sub><sup>\*</sup> + ¬IΣ<sub>1</sub><sup>0</sup> ⊢ ψ.
(ii) if ψ is Π<sub>1</sub><sup>1</sup>-conservative over RCA<sub>0</sub><sup>\*</sup>, then WKL<sub>0</sub><sup>\*</sup> + ¬IΣ<sub>1</sub><sup>0</sup> ⊢ ψ.

In contrast, the set of  $\Pi_2^1$  sentences  $\psi$  that are  $\Pi_1^1$ -conservative over RCA<sub>0</sub> is  $\Pi_2$ -complete. [Towsner 2015]

It also contains some combinatorially natural principles that do not follow from  $\mathsf{WKL}_0$ , such as the cohesive set principle COH:

"for every family  $\{R_x : x \in \mathbb{N}\}$  of subsets of  $\mathbb{N}$ , there exists infinite  $C \subseteq \mathbb{N}$  s.t. for each x, either  $\forall^{\infty}z \in C (z \in R_x)$  or  $\forall^{\infty}z \in C (z \notin R_x)$ ".

## Consequences for WKL<sub>0</sub><sup>\*</sup>: failure of low basis

#### Corollary

If  $(M, \mathcal{X}) \models RCA_0^*$ , and  $A \in \mathcal{X}$  is such that  $\neg I\Sigma_1^A$  holds, then there is a computable in A infinite 0-1 tree T such that no model  $(M, \mathcal{Y}) \models RCA_0^*$  contains any infinite path through T that is arithmetically definable in A.

- This is a failure of the low basis theorem: T is Δ<sub>1</sub>(A), but has not just no low Δ<sub>2</sub>(A) path, but even no arithmetically-in-A definable one, at least one contained in a model of RCA<sub>0</sub><sup>\*</sup>.
- In contrast, the low basis theorem is provable in RCA<sub>0</sub>. [Hájek-Kučera 1989].

## Back over RCA<sub>0</sub>

A major problem in reverse mathematics: describe the  $\Pi_1^1$  consequences of RCA<sub>0</sub> + RT<sub>2</sub><sup>2</sup>. (Here RT<sub>2</sub><sup>2</sup> is Ramsey's Thm for pairs and two colours.)

 $RCA_0 + RT_2^2$  is  $\Pi_1^1$ -conservative over  $I\Sigma_2^0$  and proves  $I\Delta_2^0$ . So, it remains to characterize its behaviour over  $I\Delta_2^0 + \neg I\Sigma_2^0$ .

But if  $(M, \mathcal{X}) \models RCA_0 + I\Delta_2^0 + \neg I\Sigma_2^0$ , then  $(M, \Delta_2^0 \text{-Def}(M, \mathcal{X})) \models RCA_0^0 + \neg I\Sigma_1^0!$ 

Is there a neat statement ensuring that  $\Delta_2^0$ -Def satisfies WKL?

## A closer look at COH

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Provably in RCA<sub>0</sub> + I\Delta_2^0, the statement
"The \Delta_2^0-definable sets satisfy WKL"
or
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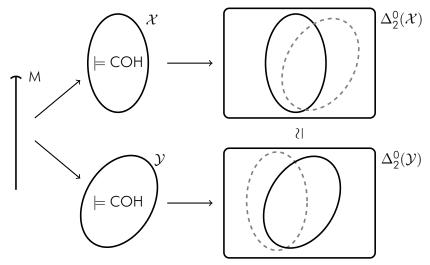
"For every set X, there is Y such that Y' has PA degree relative to X' " is equivalent to COH! [Belanger 20XX]

So, the isomorphism theorem for  $WKL_0^*$  gives:

#### Corollary

Let  $(M, \mathcal{X}), (M, \mathcal{Y})$  be countable models of  $RCA_0 + I\Delta_2^0 + COH$ . If  $(M, \mathcal{X} \cap \mathcal{Y}) \models \neg I\Sigma_2^0$ , then  $(M, \Delta_2^0 \text{-Def}(M, \mathcal{X})) \simeq (M, \Delta_2^0 \text{-Def}(M, \mathcal{Y}))$ .





## A consequence of $RT_2^2$

 $\mathrm{RT}_2^2$  says: "For every f:  $[\mathbb{N}]^2 \to 2$  there is a homogeneous set H".

Consider the following sentence  $\gamma$ :

"For every Z, if  $\neg I\Sigma_2^Z$ , then for every f:  $[\mathbb{N}]^2 \rightarrow 2$  with  $f \leq_T Z$ and every set Y such that Y' has PA degree relative to Z', there is a  $\Delta_2^0$ -set  $\widetilde{H}$  homogeneous for f s.t.  $(\widetilde{H} \oplus Z)' \leq_T Y'$ ."

(For those who care:  $\gamma$  says that if  $I\Sigma_2^0$  fails then the first-jump control argument of [CJS 2001] has to work for adding homogeneous sets for 2-colourings of pairs.)

 $\triangleright \gamma \text{ is } \Pi_1^1, \text{ in fact } \forall \Pi_5^0.$ 

►  $RCA_0 + RT_2^2 \vdash \gamma$ . (Clear over  $I\Sigma_2^0$ . Over  $I\Delta_2^0 + \neg I\Sigma_2^0$ , argue using  $RCA_0 + RT_2^2 \vdash COH$  and the iso thm for COH.)

► If  $RCA_0 + I\Delta_2^0 \vdash \gamma$ , then  $RCA_0 + RT_2^2$  is  $\Pi_1^1$ -conservative over  $I\Delta_2^0$ .

# Characterizing conservativity of $RT_2^2$

Corollary  $RCA_0 + RT_2^2$  is  $\Pi_1^1$ -conservative over  $RCA_0 + I\Delta_2^0$  iff  $RCA_0 + RT_2^2$  is  $\forall \Pi_5^0$ -conservative over  $RCA_0 + I\Delta_2^0$ .

Note:

- ►  $RCA_0 + RT_2^2$  is  $\forall \Pi_3^0$ -conservative over  $RCA_0 + I\Delta_2^0$ . [PY 2018]
- $\begin{array}{l} \blacktriangleright \ RCA_0 + RT_2^2 \text{ is } \forall \Pi_4^0 \text{-conservative over} \\ RCA_0 + I\Delta_2^0 + \{WO(\omega), WO(\omega^\omega), \ldots\}. \ \text{(Essentially [CSY 2017].)} \end{array}$
- ▶  $RCA_0^* + RT_2^2$  is  $\forall \Pi_3^0$  but not  $\forall \Pi_4^0$ -conservative over  $RCA_0^*$ . [KKY 20XX]

## Another result on conservativity

#### Theorem (Towsner 2015)

For each n, the set of  $\Pi_2^1$  sentences  $\psi$  that are  $\Pi_1^1$ -conservative over RCA<sub>0</sub> +  $I\Sigma_n^0$  is  $\Pi_2$ -complete.

Towsner asked whether this also works for  $I\Delta_n^0$  in place of  $I\Sigma_n^0$ . Much of our analysis of  $RT_2^2$  carries over to any  $\Pi_2^1$  sentence, giving:

#### Corollary (of the isomorphism thm)

For each n, the set of  $\Pi_2^1$  sentences  $\psi$  that are  $\Pi_1^1$ -conservative over  $\text{RCA}_0^* + |\Delta_n^0 + \neg |\Sigma_n^0$  is c.e.

However, a completely different argument shows:

#### Theorem

For each n, the set of  $\Pi_2^1$  sentences  $\psi$  that are  $\Pi_1^1$ -conservative over RCA<sub>0</sub><sup>\*</sup> +  $I\Delta_n^0$  is  $\Pi_2$ -complete.