

# Three open questions

André Nies

The University of Auckland

OW computability meeting Feb 2012

## The ideal of $K$ -trivial Turing degrees

- ▶ Each  $K$ -trivial is Turing below a c.e.  $K$ -trivial.
- ▶ Lots of characterizations of the  $K$ -trivial, but really not much known that is purely degree theoretic (not mentioning randomness).
- ▶ Pretty much all we know:  
each  $K$ -trivial is superlow, there is cone avoiding below a low c.e. degree, the  $\text{low}_2$  c.e. upper bound, and the low Turing cover of Kučera/Slaman (2005)

### Question (Miller, N 2006)

*Does the ideal  $\mathbf{K}$  of  $K$ -trivial Turing degrees have an exact pair  $\mathbf{a}, \mathbf{b}$  in the c.e. degrees?*

*That is, can we make  $\mathbf{K} = [\mathbf{0}, \mathbf{a}] \cap [\mathbf{0}, \mathbf{b}]$ ?*

Is the ideal definable without parameters in any Turing degree structure (C.e.,  $\Delta_2^0$ , all)?

# $\Pi_1^1$ randomness

## Definition

We say a set  $Z \subseteq \mathbb{N}$  is  $\Pi_1^1$ -random if  $Z$  is in no null  $\Pi_1^1$  class.

- ▶ There is a universal test, that is, a largest  $\Pi_1^1$  null class  $\mathcal{Q}$ , by Kechris (1975); Hjorth, N 2007.
- ▶ A  $\Pi_1^1$ -random can be  $\leq_T$  Kleenes  $\mathcal{O}$  and hyperlow ( $\omega_1^X = \omega_1^{CK}$  by Gandy basis theorem).
- ▶ Two characterizations via randomness enhancement known (weaker randomness notion + lowness property)

## Lowness for $\Pi_1^1$ randomness

We say an oracle  $A$  is low for  $\Pi_1^1$  randomness if  $\mathcal{Q}^A = \mathcal{Q}$ .

Question (Miller, N 2006; Hjorth/N 2007)

*Is such an oracle necessarily hyperarithmetical?*

We know: low for  $\Delta_1^1$  randomness =  $\Delta_1^1$  traceable (Chong, N, Yu 2008);  
low for  $\Pi_1^1$  ML-rdness = hyperarithmetical (Hj/N 2007)

(Question discussed with Harrington, Slaman, S. Friedman and co-authors on these papers.)

We could also study  $\Pi_1^1$  random hyperdegrees.

# Differentiability

General direction: in  $n$ -cube  $[0, 1]^n$ , characterize algorithmic randomness notions by differentiability of effective functions.

For instance:

## Question

$z \in [0, 1]^n$  is computably random  $\Leftrightarrow$  each computable Lipschitz function  $f: [0, 1]^n \rightarrow \mathbb{R}$  is differentiable at  $z$ ?

This is true for  $n = 1$  by work of Freer, Kjos-Hanssen and N.