## **Balancing Randomness**

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(Joint with Laurent Bienvenu, Noam Greenberg, Antonín Kučera and André Nies)

7 February, 2012

Theorem (Hirschfeldt, Nies & Stephan) Every c.e. set below an incomplete MLR is K-trivial.

Question Is every K-trivial computable from an incomplete MLR?

Theorem (Kučera & Slaman)

There is a low degree which bounds all K-trivials.

Question

Is there a low MLR which computes every K-trivial?

# Theorem There is a K-trivial A such that all MLR $X >_T A$ are super-high $(X' \ge_{tt} \emptyset'')$ .

Suppose f is a nondecreasing function. Then the derivative exists almost everywhere. So if we effectivize f, some form of randomness implies f'(z) exists.

Suppose *M* is a martingale. Then  $\lim_{n} M(X \upharpoonright_{n})$  exists almost everywhere. So if we effective *M*, some form of randomness implies  $\lim_{n} M(X \upharpoonright_{n})$  exists.

We effectivize nondecreasing functions with interval c.e. functions:

## Definition

An increasing function f is *interval c.e.* if for all rationals a < b, f(b) - f(a) is uniformly left c.e..

Example: the variation of a computable function is interval c.e..

We effectivize martingales with left c.e. martingales. (Not supermartingales.)

#### Question

For which reals Z does f'(Z) exist for all non decreasing, interval c.e. functions f?

#### Question

For which reals Z does every left c.e. martingale converge?

## Definition

A balanced test is a nested sequence of  $\Sigma_1^0$  classes  $\langle \mathcal{U}_n \rangle$  with  $\lambda(\mathcal{U}_n) \leq 2^{-n}$ , and the indices of the  $\mathcal{U}_n$  are given by a computable approximation function which changes at most  $2^n$  times on input n.

A real Z passes  $\langle \mathcal{U}_n \rangle$  if  $Z \notin \bigcap \mathcal{U}_n$ .

A real Z is balanced random if it passes every balanced test.

Weak Demuth random  $\Rightarrow$  balanced random  $\Rightarrow$  difference random (all implications proper)

#### Theorem

If Z is balanced random, then f'(Z) exists for all nondecreasing, interval c.e. functions f.

Idea of proof (for martingale version): Fix rationals  $\alpha < \beta$  such that  $\exists^{\infty} n \ M(Z \upharpoonright_n) < \alpha$ ,  $\exists^{\infty} n \ M(Z \upharpoonright_n) > \beta$ .

An *upcrossing* is a pair of strings  $\sigma \prec \tau$  with  $M(\sigma) \leq \alpha$ ,  $M(\tau) > \beta$ .

Create  $U_n$  by searching for upcrossings.

Every time  $M(\langle \rangle)$  passes a multiple of  $2^{-n}$ , change the version of  $U_n$ .

#### Definition

For a  $\Pi_1^0$  class  $C \subseteq \mathbb{R}$  and an interval  $I \subseteq \mathbb{R}$ , the density of C in I is  $\rho_C(I) = \lambda(I \cap C)/\lambda(I)$ .

For  $Z \in \mathbb{R}$ , the density of Z in C is  $\rho_{\mathcal{C}}(Z) = \liminf_{\substack{|I| \to 0 \\ Z \in I}} \rho_{\mathcal{C}}(I)$ .

(Note that  $\rho_{\mathcal{C}}([\sigma])$  is a right c.e. martingale.)

By Lebesgue density, almost every  $Z \in C$  has density 1 in C.

Theorem (Bienvenu, Hölzl, Miller, Nies) If Z is MLR and not LR-hard, then for every  $\Pi_1^0 \mathcal{C} \ni Z$ ,  $\rho_{\mathcal{C}}(Z) = 1$ . If f'(Z) exists for all non-decreasing, interval c.e. functions f, then  $\rho_{\mathcal{C}}(Z) = 1$  for all  $\Pi_1^0 \mathcal{C} \ni Z$ .

So if Z is balanced random, then  $\rho_{\mathcal{C}}(Z) = 1$  for all  $\Pi_1^0 \mathcal{C} \ni Z$ .

Does this tell us anything new? Is there an LR-hard balanced random?

If f'(Z) exists for all non-decreasing, interval c.e. functions f, then  $\rho_{\mathcal{C}}(Z) = 1$  for all  $\Pi_1^0 \mathcal{C} \ni Z$ .

So if Z is balanced random, then  $\rho_{\mathcal{C}}(Z) = 1$  for all  $\Pi_1^0 \mathcal{C} \ni Z$ .

Does this tell us anything new? Is there an LR-hard balanced random?

We couldn't build one.

As Carl mentioned, when you have problems, define them away.

The problem was that  $U_n$  might use a lot of its changes, while  $U_{n+1}$  still has most of its changes left.

## Definition

An *Oberwolfach test* is a balanced test such that every time  $U_{n+1}$  changes twice,  $U_n$  changes at least once.

A real is Oberwolfach random if it passes every Oberwolfach test.

 $\mathsf{Balanced} \ \mathsf{random} \Rightarrow \mathsf{Oberwolfach} \ \mathsf{random} \Rightarrow \mathsf{difference} \ \mathsf{random}$ 

#### Theorem

If Z is Oberwolfach random, then f'(Z) exists for all nondecreasing interval c.e. functions f.

Does this fix the problem? No.

But it gets us some nice properties:

#### Theorem

If Z is MLR but not Oberwolfach random, then Z is h-JT-hard for any computable order h with  $\sum_{x} \frac{1}{h(x)} < \infty$ , and hence Z is super-high.

#### Theorem

There is a K-trivial c.e. set A which is not below any Oberwolfach random.

Proof sketch (for balanced random):

Let  $\mathcal{E} = \{ Y \mid \exists x \in A[\Theta^Y(x) = 0] \}.$ 

Beginning at stage  $s_0$ , pick a  $v \notin A$ , and let  $\mathcal{U}_n = \{ Y \notin \mathcal{E}_{s_0} \mid A \upharpoonright_{\nu+1} \prec \Theta^Y \}.$ 

When  $\lambda(\mathcal{U}_n)$  reaches  $2^{-n}$ ...

- 1. if  $\lambda(\mathcal{U}_n \cap \mathcal{E}_s) > 2^{-(n+1)}$ , choose a large v and start over.
- 2. if  $c_{\mathcal{K}}(v,s) > 2^{-(n+1)}$ , choose a large v and start over.
- 3. otherwise, enumerate v into A, choose a large v and start over.

Putting these two together,

Corollary

There is a K-trivial A such that all  $X >_T A$  which are MLR are super-high.

Also,

Theorem (Figueira, Hirschfeldt, Miller, Ng, Nies) If X is not balanced random, then every  $Y \in MLR^X$  is balanced random.

Corollary If  $X \oplus Y$  is MLR, then at least one of X, Y is balanced random.

Corollary

There is a K-trivial which is not below both halves of a MLR.

## Question

*Can we separate Oberwolfach randomness from difference randomness?*