

Who asked *us*?

How the theory of computing answers questions
that weren't about computing

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Computability Theory
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Algorithmic Information (Kolmogorov Complexity)

The **Kolmogorov complexity** of a string $x \in \{0, 1\}^*$ is

$$K(x) = \min \{ |\pi| \mid \pi \in \{0, 1\}^* \text{ and } U(\pi) = x \},$$

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- $K(x) \leq |x| + o(|x|)$.
- x is “random” if $K(x) \approx |x|$.
- Routine coding extends this to $K(x)$ for $x \in \mathbb{N}$, $x \in \mathbb{Q}$, $x \in \mathbb{Q}^n$, etc.

Dimensions of Points

Work in Euclidean space \mathbb{R}^n .

The **Kolmogorov complexity** of a set $E \subseteq \mathbb{Q}^n$ is

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(Shen and Vereschagin 2002)

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Note that

$$E \subseteq F \Rightarrow K(E) \geq K(F).$$

Let $x \in \mathbb{R}^n$ and $r \in \mathbb{N}$. The **Kolmogorov complexity** of x at **precision** r is

$$K_r(x) = K(B_{2^{-r}}(x)),$$

i.e., the number of bits required to specify **some** rational point $q \in \mathbb{Q}^n$ such that $|q - x| \leq 2^{-r}$.

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Old fact (J. Lutz '00 + Hitchcock '03). If $E \subseteq \mathbb{R}^n$ is a union of Π_1^0 sets, then

$$\dim_H(E) = \sup_{x \in E} \dim(x).$$

classical Hausdorff
(fractal) dimension

dimensions of
individual points

\therefore Dimensions of points are geometrically meaningful.

Dimensions of Points

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$\text{Dim}(x)$ is the “ Σ_1^0 version” of \dim_P . (packing dimension)

Theorem (J. Lutz and N. Lutz, STACS '17)

For every $E \subseteq \mathbb{R}^n$,

$$\dim_H(E) = \min_{A \subseteq \mathbb{N}} \sup_{x \in E} \dim^A(x).$$

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or, if you're lucky, that

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Theorem (J. Lutz and N. Lutz, STACS '17)

For *every* $E \subseteq \mathbb{R}^n$,

$$\dim_P(E) = \min_{A \subseteq \mathbb{N}} \sup_{x \in E} \text{Dim}^A(x).$$

A **Keakeya set** in \mathbb{R}^n is a set $K \subseteq \mathbb{R}^n$ that contains a unit segment in every direction.

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Keakeya Conjecture. Every Keakeya set in \mathbb{R}^n has Hausdorff dimension n .

- An important open problem for $n \geq 3$.

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It doesn't seem to help in \mathbb{R}^n for $n \geq 3$.

Dimensions of Points on $y = mx + b$

Question (J. Lutz, early 2000s). Is there a line $y = mx + b$ on which **every** point has dimension 1?

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$$\dim(x, mx + b) \geq \dim^{m,b}(x) + \min\{\dim(m, b), \dim^{m,b}(x)\}.$$

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Corollary. For every $m, b \in \mathbb{R}$ there exist $x_1, x_2 \in \mathbb{R}$ such that

$$\dim(x_1, mx_1 + b) - \dim(x_2, mx_2 + b) \geq 1.$$

\therefore The answer to the above question is “No!”

Generalized Furstenberg sets

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For $\alpha \in (0, 1]$, a set $E \subseteq \mathbb{R}^2$ is **α -Furstenberg** if, for every $e \in S^1$ (= the unit circle in \mathbb{R}^2), there is a line \mathcal{L}_e in direction e such that $\dim_H(\mathcal{L}_e \cap E) \geq \alpha$.

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Definition (Molter and Rela 2012)

For $\alpha, \beta \in (0, 1]$, a set $E \subseteq \mathbb{R}^2$ is **(α, β) -generalized Furstenberg** if there is a set $J \subseteq S^1$ such that $\dim_H(J) \geq \beta$ and, for every $e \in J$, there is a line \mathcal{L}_e in direction e such that $\dim_H(\mathcal{L}_e \cap E) \geq \alpha$.

Theorem (probably Furstenberg and Katznelson)

For $\alpha \in (0, 1]$, every α -Furstenberg set $E \subseteq \mathbb{R}^2$ satisfies

$$\dim_H(E) \geq \alpha + \max \left\{ \frac{1}{2}, \alpha \right\}.$$

Note that Davies's theorem follows from the case $\alpha = 1$.

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For $\alpha, \beta \in (0, 1]$, every (α, β) -generalized Furstenberg set $E \subseteq \mathbb{R}^2$ satisfies

$$\dim_H(E) \geq \alpha + \max \left\{ \frac{\beta}{2}, \alpha + \beta - 1 \right\}.$$

Note that the previous theorem is the case $\beta = 1$.

Theorem (N. Lutz and D. Stull, TAMC '17)

For $\alpha, \beta \in (0, 1]$, every (α, β) -generalized Furstenberg set $E \subseteq \mathbb{R}^2$ satisfies

$$\dim_H(E) \geq \alpha + \min\{\beta, \alpha\}.$$

Note that this improves on the theorem of Molter and Rela exactly when $\alpha < 1$, $\beta < 1$, and $\beta < 2\alpha$. Hence it doesn't improve the bound on α -Furstenberg sets.

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It is the **first use** of **algorithmic** fractal dimensions to prove a **new** theorem in **classical** fractal geometry!

Intersections and Products of Fractals

The following are fundamental, nontrivial, textbook theorems of fractal geometry.

Product Formula (Marstrand 1954). For all sets $E \subseteq \mathbb{R}^m$ and $F \subseteq \mathbb{R}^n$,

$$\dim_H(E \times F) \geq \dim_H(E) + \dim_H(F).$$

Intersection Formula (Kahane 1986; Mattila 1984, 1985). For all **Borel** sets $E, F \subseteq \mathbb{R}^n$ and almost every $z \in \mathbb{R}^n$,

$$\dim_H(E \cap (F + z)) \leq \max\{0, \dim_H(E \times F) - n\}.$$

Note: The product formula was known earlier with extra assumptions on E and F . Marstrand deployed nontrivial machinery to prove it for arbitrary sets.

Textbooks usually just prove it for Borel sets.

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This paper also uses a similar method to give a **much** simpler proof of the general Product Formula, along with analogous results for packing dimension.

Classical fractal geometry has a pointwise notion of dimension.

An **outer measure** on \mathbb{R}^n is a function $\nu : \mathcal{P}(\mathbb{R}^n) \rightarrow [0, \infty]$ satisfying

- $\nu(\emptyset) = 0$,
- $E \subseteq F \Rightarrow \nu(E) \leq \nu(F)$, and
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An outer measure ν on \mathbb{R}^n is

- **finite** if $\nu(\mathbb{R}^n) < \infty$, and
- **locally finite** if every $x \in \mathbb{R}^n$ has a neighborhood N with $\nu(N) < \infty$.

Definition

Let ν be a locally finite outer measure on \mathbb{R}^n , and let $x \in \mathbb{R}^n$. The **lower** and **upper pointwise dimensions** of ν at x are

$$\dim_{\nu}(x) = \liminf_{r \rightarrow \infty} \frac{\log \frac{1}{\nu(B_{2^{-r}}(x))}}{r}$$

and

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Are these in any way related to the algorithmic dimensions $\dim(x)$ and $\text{Dim}(x)$?

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For each $E \subseteq \mathbb{R}^n$, let

$$\kappa(E) = 2^{-K(E)}.$$

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2. For all $x \in \mathbb{R}^n$, $\dim(x) = \dim_{\kappa}(x)$ and $\text{Dim}(x) = \text{Dim}_{\kappa}(x)$.
3. This relativizes and interacts informatively with the Point-to-Set Principle.

Thank you!