

# On definability of c.e. degrees in the 2-c.e. degree structures

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## Definitions and notations

All sets are subsets of  $\omega = \{0, 1, 2, \dots\}$ . Thus, let  $A, B \subset \omega$ .

- $A \leq_T B$  if there is an algorithm that allows to answer the questions “ $x \in A$ ?”, using  $B$  as an oracle.
- $A \leq_m B$  if there is a computable function  $f$  such that  $x \in A \iff f(x) \in B$ .
- Clearly,  $A \leq_m B \implies A \leq_T B$ .

## Definitions and notations

- If  $A \leq_r B$  and  $B \leq_r A$  then  $A \equiv_r B$ .
- Let  $\text{deg}_r(A) = \{B \mid A \equiv_r B\}$ .
- Here,  $r \in \{m, T\}$ .

# The Coopers theorem

## Theorem (Cooper, 1971)

There is a 2-c.e. set with proper 2-c.e. Turing degree.

**Remark.** A 2-c.e. degree is proper if it doesn't contain a c.e. set. In particular, the constructed set has proper 2-c.e.  $m$ -degree.

As a corollary, the universes for c.e. and 2-c.e. degree structures are different. And clearly, c.e. degrees form a substructure in the corresponding 2-c.e. degrees.

## Motivations and goals

- To investigate the 2-c.e. degree structures.
- To investigate model-theoretic properties of c.e. and 2-c.e. degrees (in different settings).
- To study relationship between c.e. and properly 2-c.e. degrees (in different settings).

## Motivations and goals

### Open question (Cooper, 2002; Arslanov, 2009)

Is the class of c.e. Turing degrees definable in the partial ordering of 2-c.e. Turing degrees?

Related questions:

- The same questions for  $m$ -degrees.
- A weaker version of the question involving parameters.
- A weaker version of the question involving additional predicates.
- The case of low c.e. and 2-c.e. degrees.

# Definability

Let  $\mathcal{A}$  be a structure, and  $B$  be a subset of  $|A|$ .

## Definition

The class  $B$  is definable in  $\mathcal{A}$  if there exists a formula  $\varphi(x)$  of the first order language such that for all  $a \in |A|$  it holds

$$\mathcal{A} \models \varphi(a) \Leftrightarrow a \in B$$

- As  $\mathcal{A}$  we consider  $(\mathbf{D}, \leq)$ , where  $\mathbf{D}$  is the corresponding 2-c.e. degrees, and  $\leq$  is induced by the same reducibility.
- As  $B$  we consider  $\mathbf{R}$ , the c.e. degrees.
- In  $\varphi$ , there can be additional fixed variables  $c_1, c_2, \dots \in |A|$  called parameters.

# Section 1

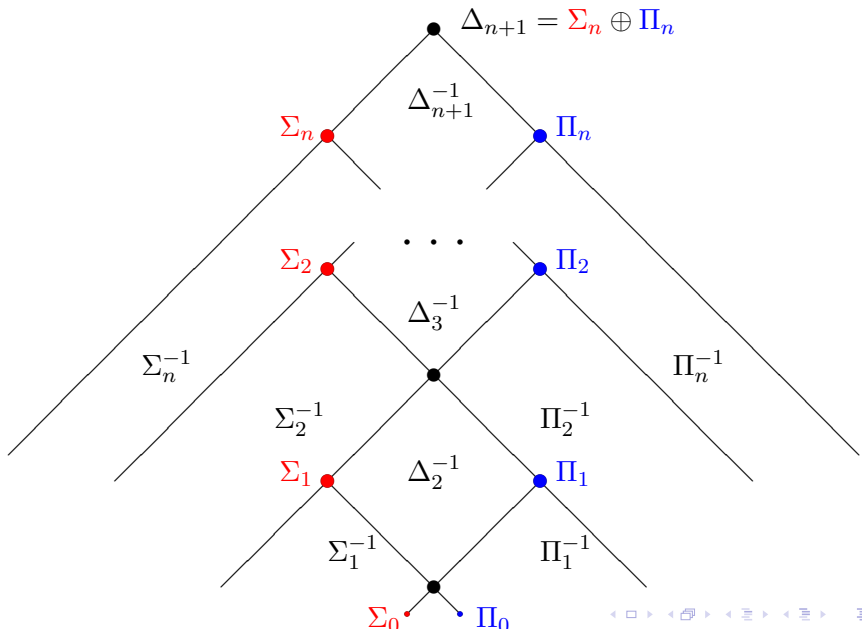
Definability for  $m$ -degrees



## A brief history

- $m$ -degrees were actively studied since 1970 (by Degtev, Denisov, Ershov, Nies, Lachlan, Selivanov, etc. )
- The most attention was received by c.e.  $m$ -degrees and by all  $m$ -degrees.
- In general the structures of  $m$ -degrees found out to have many good properties, in particular, much better than the structures of  $T$ -degrees.
- For example,  $\Sigma_n^{-1}$   $m$ -degrees have the greatest (universal) element (by Ershov),  $\Delta_n^{-1}$   $m$ -degrees have the greatest element (by Selivanov), for any fixed  $n > 0$ .

# Picture



## Facts and folklore

- c.e.  $m$ -degrees form an ideal in 2-c.e.  $m$ -degrees
- c.e. and co-c.e.  $m$ -degrees are isomorphic.
- c.e. and 2-c.e.  $m$ -degrees form a distributive upper semilattice (by Ershov, Lachlan, Selivanov). The same holds for c.e.  $wtt$ -degrees, but doesn't hold for 2-c.e.  $wtt$ -degrees, and for c.e. and 2-c.e. Turing degrees.
- The greatest c.e.  $m$ -degree is not splittable (by Lachlan), thus  $\Delta$ - and  $\Sigma$ -( $\Pi$ -) levels are not elementarily equivalent. The result has a direct generalization to 2-c.e.  $m$ -degrees.

## Facts and folklore

- Given 2-c.e. set  $A = A_0 - A_1$ , let  $A_0 = \text{rng}(f)$  for some computable 1-1  $f$ , then  $L(A) = f^{-1}(A_1)$  is Lachlan's set for  $A$ .
- $L(A)$  is c.e.
- $\overline{L(A)} \leq_m A$
- If  $L(A)$  is c.e. then  $A$  is 2-c.e., and if  $L(A)$  is computable then  $A$  is c.e.
- Then below any proper 2-c.e.  $m$ -degree there exists a noncomputable co-c.e.  $m$ -degree.

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## Elementary difference

### Theorem (Ershov and Lavrov, 1973)

*Given noncomplete c.e. set  $B$  there exists a c.e. set  $A \not\leq_m B$  which is minimal*

### Corollary

*c.e. and 2-c.e.  $m$ -degrees are not elementarily equivalent*

Note that for we can take  $U_{\Delta_2}$  as  $B$ , then any set  $A \not\leq_m B$  would be proper 2-c.e. and have a noncomputable element below it.

*Remark.* The theorem was proved for much more general case and for the c.e. setting. For details see [Erhov Yu.L., Lavrov I.A. Upper semilattice  $L(S)$ , Algebra i Logica, 1973, Vol.12, No.2, P.167-189]

# The main structural theorems

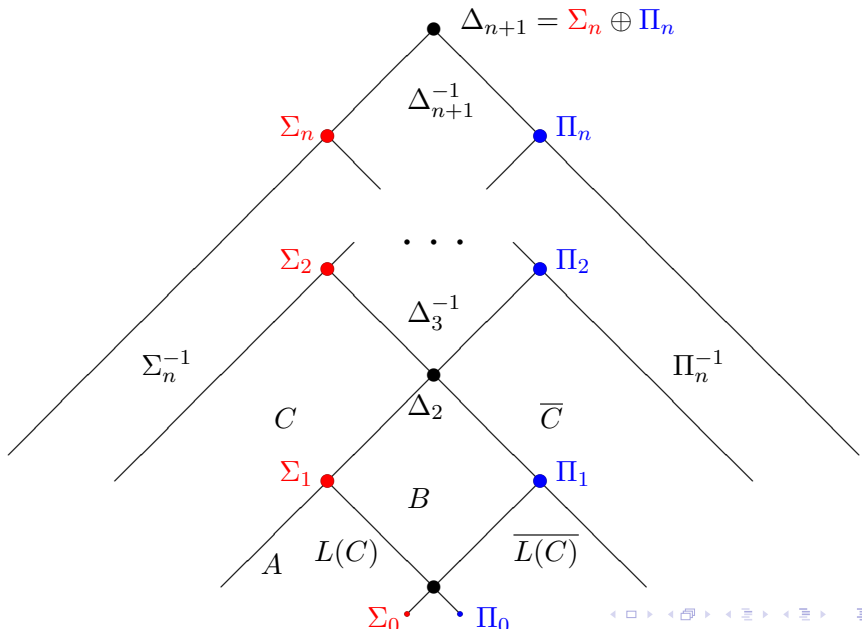
## Theorem 1 (Ng, Yamaleev)

Given  $k, n > 0$ , given any  $\Sigma_k^{-1}$  set  $B$  such that  $U_{\Sigma_n} \not\leq_m B$ , there exists a  $\Sigma_n^{-1}$  set  $A \not\leq_m B$  such that for any  $W <_m A$  it holds that  $W \leq_m U_{\Delta_n}$ .

- noncomplete  $B \iff U_{\Sigma_n} \not\leq_m B$
- minimal  $A \iff A$  is minimal cover for  $U_{\Delta_n}$
- c.e.  $A \not\leq_m B \iff \Sigma_n^{-1}$  set  $A \not\leq_m B$



# Picture



## The main structural theorems

### Corollary (Ng, Yamaleev)

Given  $k, n > 0$ , given any  $\Sigma_k^{-1}$  set  $B$  such that  $U_{\Delta_{n+1}} \not\leq_m B$ , there exists a  $\Delta_{n+1}^{-1}$  set  $A \not\leq_m B$  such that for any  $W <_m A$  it holds that  $W \leq_m U_{\Delta_n}$ .

### Theorem 2 (Ng, Yamaleev)

Given  $n > 0$ , there exists a set  $A$  of properly  $\Sigma_{n+1}^{-1}$  degree such that for any  $W \in \Sigma_n^{-1}$  if  $W \leq_m A$  then  $W \leq_m U_{\Delta_n}$ .

## The main structural theorems, 2-c.e. setting

### Corollary (Ershov, Lavrov, 1973)

Given noncomplete (in  $\Delta_2^{-1}$   $m$ -degrees) set  $B$  there exists a 2-c.e. set  $A \not\leq_m B$  such that  $A$  has a minimal  $m$ -degree (moreover, it will be either c.e. or co-c.e.)

### Corollary (from Theorem 2 (Ng, Yamaleev))

There exists a set  $A$  of properly 2-c.e.  $m$ -degree such that for any c.e.  $W$  if  $W \leq_m A$  then  $W$  is computable (i.e.,  $A$  form minimal pair with the greatest c.e. degree).

## Intuitive description

- The first part says we can build minimal  $m$ -degrees avoiding arbitrary (noncomplete) lower cones.
- The second part says that for all c.e.  $m$ -degrees we can find a half minimal pair in the 2-c.e.  $m$ -degrees.
- Note also that we cannot do it for co-c.e.  $m$ -degrees using a unique 2-c.e. degree.

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## The corollaries

- The degree structures of c.e. and 2-c.e.  $m$ -degrees are not elementarily equivalent (and it works for all higher levels).
- The  $m$ -degree of universal  $\Delta_2^{-1}$ -set is definable in 2-c.e.  $m$ -degrees.
- The complementary Theorem 2 allows to distinguish the greatest c.e. from the greatest co-c.e.  $m$ -degree.

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## Definability of c.e. in 2-c.e.

- $\theta(x) := \forall b [x \not\leq b \Rightarrow \exists a (a \not\leq b \wedge \forall w [w < a \Rightarrow w \leq 0])]$
- $\psi(x) := \theta(x) \wedge \forall z [x < z \Rightarrow \neg\theta(z)],$
- Thus,  $\psi(x)$  is true in  $\Sigma_2^{-1}$  iff  $x = U_{\Delta_2}$
  
- $\varphi(x, y) := \exists u \psi(u) \wedge x \cup y = u \wedge [\forall x_1 \forall y_1 (x_1 < x \Rightarrow x_1 \cup y < u) \wedge (y_1 < y \Rightarrow x \cup y_1 < u)]$
- Thus,  $\varphi(x, y)$  defines the pair of  $U_{\Sigma_1}$  and  $U_{\Pi_1}$  but cannot distinguish them.
  
- $\varphi^\Sigma(x) := \exists y (\varphi(x, y) \wedge \exists z \forall w [z \not\leq x \cup y \wedge w < z \wedge w \leq x \Rightarrow w \leq 0])$

## Complexity of the formulas

- Elementarily difference of c.e. and 2-c.e.:  $\Sigma_2^0$
- Definability of c.e. in 2-c.e.:  $\Sigma_4^0$
- For higher levels the complexity grows incredibly. For instance, in  $\Sigma_n^{-1}$  we define in the following ordering:  $U_{\Delta_2}, U_{\Delta_3}, \dots, U_{\Delta_n}$ , then  $U_{\Sigma_{n-1}}, U_{\Sigma_{n-2}}, \dots, U_{\Sigma_1}$ .

## Questions

- Is  $\Sigma_1^{-1}$  level definable in the structure of  $\Sigma_\omega^{-1}$ -level?
- Could the same approach work for infinite levels? (probably with parameters)
- What to do with the limit levels?

## Section 2

A weaker definability for Turing degrees

# Approaches

to the problem of definability of c.e. Turing degrees in partial ordering of 2-c.e. Turing degrees.

Proposed by Arslanov and Yamaleev (2018)

1. Density of double bubbles
2. Nonspilliting pairs
3. Lachlan sets and degrees
4. Isolation from side

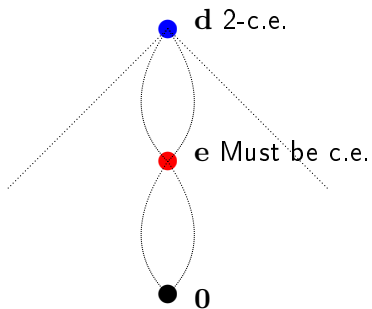
## Definable singletons

### Definition (Arslanov, Kalimullin, and Lempp, 2010)

Let  $\mathbf{e}, \mathbf{d}$  be 2-c.e. degrees such that  $\mathbf{0} < \mathbf{e} < \mathbf{d}$ . We say that these degrees form a *double bubble* (also, a *bubble pair*, *2-bubble*, *bubble*) in 2-c.e. degrees if any 2-c.e. degree  $\mathbf{u} < \mathbf{d}$  is comparable with  $\mathbf{e}$ . Also, we say that  $\mathbf{d}$  is the top of bubble, and  $\mathbf{e}$  is the middle of bubble. By default, we consider bubbles in 2-c.e. degrees.

- The degree  $\mathbf{e}$  must be c.e.
- The degree  $\mathbf{d}$  is an exact 2-c.e. degree.
- The degree  $\mathbf{d}$  is not splittable avoiding upper cone of  $\mathbf{e}$ .

## Approach 1. The picture.



## Approach 1. The idea.

- To show that between any two c.e. degrees we can find a degree  $e$ .
- Then any c.e. degree has a splitting where the both parts are middles of bubbles.
- Such splitting doesn't exist for properly 2-c.e. degrees.



## Approach 1. The results.

- [Liu, Wu, Yamaleev, 2015] The exact 2-c.e. degrees are downward dense.
- [Andrews, Kuyper, Lempp, Soskova, Yamaleev, 2017] There exists a nonzero c.e. degree such that no double bubble can be found below it.
- Conjecture [Arslanov, Yamaleev, 2018] The middles of double bubbles can be found below any nonzero c.e. degree, moreover it can be combined with lower cone avoidance.

## Approach 1. Conclusion.

- Definable middle of bubbles with fairly “easy” construction.
- Even if we cannot prove the density. The middles of bubbles is still a reliable class of c.e. degrees. And can be combined with downward density and cone avoidance.
- Can a middle of bubble be constructed above any low or superlow c.e. degree?

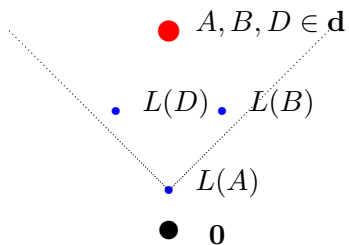
## Approach 3. Idea.

- To use  $L(D)$  which reflects to enumerability properties of a 2-c.e. set  $D$ . Then consider a collection of  $L(B)$  such that  $B \equiv_T D$ .
- Make a connection between the associated degrees  $L(B)$  and the degree of  $D$ .
- The good case is when for each properly 2-c.e. degree of  $B$  the collection of the degrees of  $L(B)$  is bounded from below by some nonzero c.e. degree.

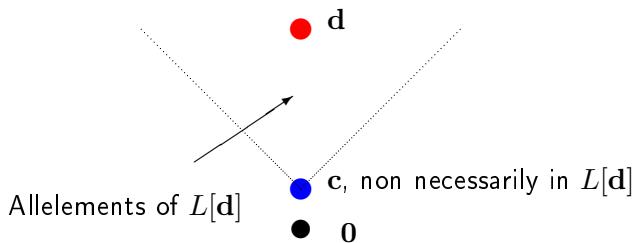
## Approach 3. Results.

- Series results by Ishmukhametov [1999,2000] and by Fang, Liu, Wu, Yamaleev [2013-2019] showed that different distributions for  $L(B)$  are possible.
- In particular, there is a properly 2-c.e. degree with unbounded collection of its associated degrees of  $L(B)$ .
- Also: if  $D \equiv_T B$  and have a proper 2-c.e. degree then  $L(D)$  and  $L(B)$  cannot form a minimal pair.

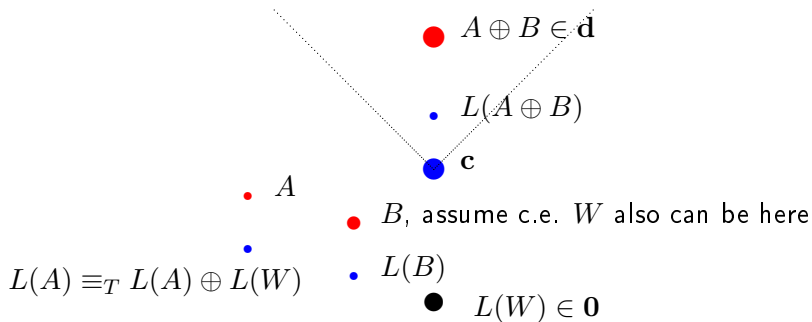
### Approach 3. Picture.



## Lower bounds for $L[\mathbf{d}]$ ?



## Approach 3. Motivation



## Lachlan degrees

- We say that  $L[\mathbf{d}] = \{\deg(L(D)) \mid D \in \mathbf{d}\}$  is a spector of Lachlan degrees for  $\mathbf{d}$ .
- Let  $R[\mathbf{d}] = \{W \mid W \text{ c.e. and } \mathbf{d} \text{ is } CEA(W)\}$ .
- Then, clearly,  $L[\mathbf{d}] \subset R[\mathbf{d}]$ .
- The following theorem allows to obtain  $R[\mathbf{d}] \subset L[\mathbf{d}]$ .

### Theorem (Arslanov, LaForte, Slaman, 1998)

Given  $\omega$ -c.e. degree  $\mathbf{d}$ , which is  $CEA(\mathbf{c})$  for some c.e. degree  $\mathbf{c}$ .  
Then there exists a 2-c.e. set  $D \in \mathbf{d}$  such that  $D$  is  $CEA(\mathbf{c})$ .  
*Moreover, the degree  $\mathbf{c}$  contains  $L(D)$ .*



## Lachlan degrees

- [Ishmukhametov, 1999]. There exists a noncomputable 2-c.e. degree  $\mathbf{d}$  such that  $L[\mathbf{d}] = [\mathbf{c}, \mathbf{b}]$  for some noncomputable c.e. degrees  $\mathbf{c}$  and  $\mathbf{b}$ . In particular, it can be  $\mathbf{c} = \mathbf{b}$ .
- [Arslanov, Kalimullin, Lempp, 2010]. There exists 2-c.e. degrees  $\mathbf{c} < \mathbf{d}$  such that they form bubble. In particular, it also holds  $L[\mathbf{d}] = \{\mathbf{c}\}$ .
- For such bubble pairs the degree  $\mathbf{c}$  is definable.
- [Ishmukhametov, 1999]. Question. Does  $L[\mathbf{d}]$  always contain a least element for any  $\mathbf{d}$ ?
- [Ishmukhametov, 2000]. There exists a 2-c.e. degree  $\mathbf{d}$  such that  $L[\mathbf{d}]$  doesn't have a least element.
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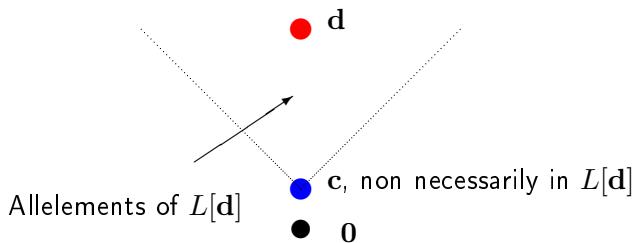
## Spectra of Lachlan degrees

- [Fang, Wu, Yamaleev, 2013]. There exists 2-c.e. degree  $\mathbf{d}$  such that  $L[\mathbf{d}]$  doesn't have a minimal element.
- [Arslanov, 2000; Fang, Liu, Wu, Yamaleev, 2015]. If  $\mathbf{d}$  is a properly 2-c.e. degree then  $L[\mathbf{d}]$  doesn't contain a minimal pair.
- **Corollary.** For any properly 2-c.e. degrees  $\mathbf{d}$  a minimal element in  $L[\mathbf{d}]$  is the least element.
- [Fang, Liu, Wu, Yamaleev]. There exists 2-c.e. degree  $\mathbf{d}$  such that  $L[\mathbf{d}]$  is not bounded from below by some nonmcomputable c.e. degree.
- [Yamaleev]. For any proper 2-c.e. degree  $\mathbf{d}$  its spector  $L[\mathbf{d}]$  differs from  $(\mathbf{0}, \mathbf{d}) \cap \mathbf{R}$ .

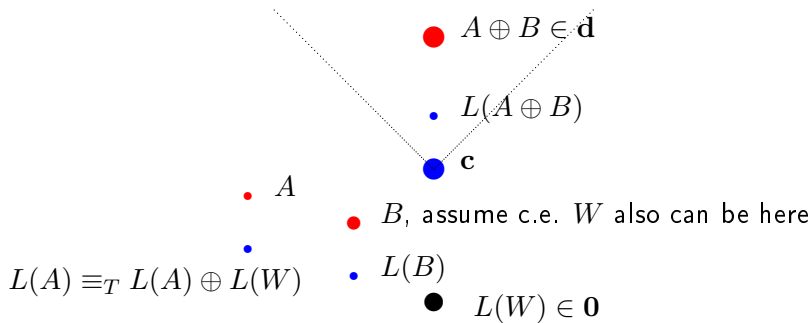
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## Lower bounds for $L[\mathbf{d}]$ ?



# Motivation



# Definability of c.e. degrees in the structures with $CEA$

Consider 2-c.e. Turing degrees  $\mathbf{D}(\leq, CEA)$ .

- [Cai, Shore, 2013]. C.e. degrees definable in  $\mathbf{D}(\leq, CEA)$  with  $\Sigma_2^0$  formula, but not with  $\Sigma_1^0$ -formula.
- [Yamaleev]. For any properly 2-c.e. degree  $\mathbf{d}$  its spector  $L[\mathbf{d}]$  differs from the interval  $(\mathbf{0}, \mathbf{d}) \cap \mathbf{R}$ .
- Corollary. C.e. degrees are definable  $\mathbf{D}(\leq, CEA)$  with a  $\Pi_1^0$ -formula.

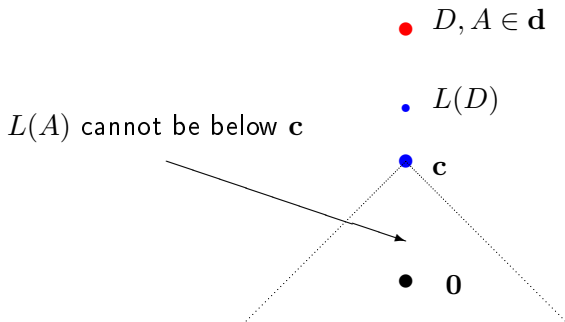


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- Corollary. C.e. degrees are definable  $\mathbf{D}(\leq, CEA)$  with a  $\Pi_1^0$ -formula.

# Definability with $\Pi_1^0$ formula



## Approach 4. Isolation from side.

- [Yang and Yu, 2006] Inapparently used isolation from side to show that c.e. degrees doesn't form a  $\Sigma_1$ -substructure of 2-c.e. degrees.
- [Cai, Slaman, and Shore, 2012] Inapparently used isolation from side to show that  $k$ -c.e. degrees doesn't form a  $\Sigma_1$ -substructure of  $n$ -c.e. degrees for all  $k < n$
- [Wu and Yamaleev, 2012] A 2-c.e. degree  $\mathbf{d}$  is *isolated from side nontrivially* if  $\mathbf{d}$  is nonisolated and there exists a c.e. degree  $\mathbf{a} \mid \mathbf{d}$  such that for all c.e. degrees  $\mathbf{w}$  if  $\mathbf{w} \leq \mathbf{d}$  then  $\mathbf{w} \leq \mathbf{a}$ .

## Approach 4. The result.

Any low properly 2-c.e. Turing degree  $\mathbf{d}$  is isolated from side

### Theorem (Yamaleev, 2019)

For any low 2-c.e. set  $D$  with a properly 2-c.e. Turing degree there exists a c.e. set  $A$  such that  $D \not\leq_T A$  and for any c.e. set  $W \leq_T D$  it holds  $W \leq_T A$ .

- The set  $A$  can be made low
- If  $D \leq_T C$  then the set  $A$  can be made below  $C$ .

## Approach 4. The consequences.

- Recall from Approach 1,  
**Conjecture [Arslanov, Yamaleev, 2018]** The middles of double bubbles can be found below any nonzero c.e. degree, moreover it can be combined with lower cone avoidance.
- In particular, for isolation from side we bound all middles of double bubbles.

### Corollary

The low c.e. degrees are definable in the partial ordering of low 2-c.e. Turing degrees

- For any low c.e. degree we can construct a definable c.e. degree below it, avoiding any lower cone
- Due to isolation from side we cannot do it for any properly low 2-c.e. degrees.

## Approach 4. The definability (with parameters).

- [Welch, 1980] There exists low c.e. degrees  $\mathbf{c}_1$  and  $\mathbf{c}_2$  such that for any c.e. degree  $\mathbf{a}$  there exists its splitting  $\mathbf{a}_1 \cup \mathbf{a}_2 = \mathbf{a}$  such that  $\mathbf{a}_i \leq \mathbf{c}_i$  for  $i = 1, 2$ .
- The parameters  $\mathbf{c}_1$  and  $\mathbf{c}_2$  are the desired ones. Lets fix them.
- Consider a c.e. degree  $\mathbf{a}$ . It has the mentioned above splitting such that the both parts are below the parameters, also those parts are not isolated from side (i.e. we always can find a definable c.e. degree below them).
- Consider a properly 2-c.e. degree  $\mathbf{d}$ . If it doesn't have a splitting below the parameters then it is clearly proper 2-c.e. Assume it has such splitting. Then at least one part must be properly 2-c.e. Then at least one part must be isolated from side (recall that  $\mathbf{c}_1$  and  $\mathbf{c}_2$  are low).

## Approach 4. Misc.

- Assume that for a given c.e. degrees  $\mathbf{a} \not\leq \mathbf{c}$  there is a middle of bubble  $\mathbf{e} < \mathbf{a}$  such that  $\mathbf{e} \not\leq \mathbf{c}$ . How to avoid the case when  $\mathbf{c}$  could be 2-c.e.?
- Then we should update isolation from side as follows: given 2-c.e. degree  $\mathbf{d}$  and c.e. degree  $\mathbf{a}$  such that  $\mathbf{a} \not\leq \mathbf{d}$ . Then there is a c.e. degree  $\mathbf{c}$  such that it covers the c.e. degrees below  $\mathbf{d}$  (can include  $\mathbf{d}$  as well) and  $\mathbf{a} \not\leq \mathbf{c}$ .

## Approach 4. Backup plans.

- For a given properly 2-c.e. degree  $\mathbf{d}$  do there exist c.e. degrees  $\mathbf{c}$  and  $\mathbf{g}$  such that one of them isolates  $\mathbf{d}$  from side?
- For a given properly 2-c.e. degree  $\mathbf{d}$  do there exist c.e. degrees  $\mathbf{c}$  and  $\mathbf{g}$  such that any c.e. degree below  $\mathbf{d}$  is either below  $\mathbf{c}$  or  $\mathbf{g}$ ?
- Note that then we obtain definable degrees which are join of two middles of bubbles? Does this class coincide with the middles of bubbles?



## Turing degrees. Conclusion

- Definability with 2 parameters.
- Definability in smaller structures (low 2-c.e. degrees).
- Definability with additional predicate  $CEA$  (at least possible level).

## Open questions

- Is any properly 2-c.e. degree isolated from side?
- Is any properly 2-c.e. degree pseudoisolated (by G. Wu, 2005)?
- Can c.e. degree be definable with 1 parameter in the partial ordering of 2-c.e. degrees?

## Section 3

The Ershov hierarchy and the CEA hierarchy

## Questions

- Given a 2-c.e. degrees. In which c.e. it can be  $CEA$ ?
- Given a c.e. degree. Which 2-c.e. degrees are  $CEA$  in it?

Open question (Soare, 1994; Arslanov, Lempp, Shore, 1996; Cooper, Li, 1998; LaForte, 2001; Arslanov, 2011)

Given low noncomputable c.e. degree  $\mathbf{c}$ , do there exists a properly 2-c.e. degree such that  $\mathbf{d}$  is  $CEA(\mathbf{c})$ ?

*(Due to the paper of Soare and Stob, 1982)*

## The CEA hierarchy

- A set  $D$  is  $CEA(C)$  if  $C \leq_T D$  and  $D$  is  $\Sigma_1^C$  ( $CEA = REA$ )
- A degree  $\mathbf{d}$  is  $CEA(\mathbf{c})$  if for some  $D \in \mathbf{d}$  and  $C \in \mathbf{c}$  we have that  $D$  is  $CEA(C)$
- A set  $A$  is  $n$ - $CEA$  if  $A$  is  $CEA(C)$  for some  $(n - 1)$ - $CEA$  set  $C$
- A degree  $\mathbf{d}$  is properly  $n$ - $CEA$  if it is  $n$ - $CEA$ , but not  $(n - 1)$ - $CEA$
- C.e. degrees are just 1- $CEA$  degrees.
- The same doesn't hold for 2-c.e. degrees.

# The CEA hierarchy

## Theorem (Soare, Stob, 1982)

Given noncomputable c.e. degree  $\mathbf{c}$ , there exists a non-c.e. degree  $\mathbf{d}$  which is  $CEA(\mathbf{c})$ .

## Theorem (Cholak, Hinman, 1994)

Given noncomputable c.e. degree  $\mathbf{c}$ , for all  $n \geq 1$  there exists a non- $n$ - $CEA$  degree  $\mathbf{d}$  which is  $CEA(\mathbf{c})$ .

**Remark.** In the first theorem  $n = 1$ , thus  $\mathbf{d}$  is 2- $CEA$ .

## Enumerability relative to low c.e.degrees

In  $\Delta_2^0$ -degrees:

**Theorem (Soare, Stob, 1982)**

Given noncomputable low c.e. degree  $\mathbf{c}$ , there exists a non-c.e. degree  $\mathbf{d}$  which is  $CEA(\mathbf{c})$

**Theorem (Arslanov, Lempp, Shore, 1996)**

There exists noncomplete c.e. degree  $\mathbf{c}$  such that any  $\Delta_2^0$ -degree, which is  $CEA(\mathbf{c})$ , must be c.e.

**Theorem (Arslanov, LaForte, Slaman, 1998)**

Given  $\omega$ -c.e. degree  $\mathbf{d}$ , which is  $CEA(\mathbf{c})$  for some c.e. degree  $\mathbf{c}$ . Then there exists a 2-c.e. set  $D \in \mathbf{d}$  such that  $D$  is  $CEA(\mathbf{c})$ .

Open question (Soare, 1994; Arslanov, Lempp, Shore, 1996; Cooper, Li, 1998; LaForte, 2001; Arslanov, 2011)

Given low noncomputable c.e. degree  $\mathbf{c}$ , do there exists a properly 2-c.e. degree such that  $\mathbf{d}$  is  $CEA(\mathbf{c})$ ?



## The negative answer

### Theorem (Arslanov, Batyrshin, Yamaleev)

There exists noncomputable low c.e. degree  $\mathbf{c}$  such that any 2-c.e. degree, which is  $CEA(\mathbf{c})$ , must be c.e.

### Corollary(Arslanov, Batyrshin, Yamaleev)

There exists noncomputable low c.e. degree  $\mathbf{c}$  such that any  $\omega$ -c.e. degree, which is  $CEA(\mathbf{c})$ , must be c.e.

## Corollaries

### Corollary (Arslanov, Batyrshin, Yamaleev)

There exists low c.e. degrees, which cannot be Lachlan degrees for properly 2-c.e. degrees.

### Corollary (Arslanov, Batyrshin, Yamaleev)

There exists low c.e. degrees  $\mathbf{b} \leq \mathbf{c}$  such that any  $\Delta_2^0$ -degree, which is  $CEA(\mathbf{b})$  and  $> \mathbf{c}$ , must be c.e.

Recall that if  $\mathbf{c}$  is superlow then non-c.e.  $CEA(\mathbf{c})$  degrees must be 2-c.e.

## Generalization and question

### Theorem (Arslanov, Batyrshin, Yamaleev)

Let  $\mathcal{U}$  be a class of  $\Delta_2^0$ -sets uniformly computable in  $\emptyset\emptyset'$ . Does there exist a low c.e. degree  $\mathbf{c}$  such that any set from  $\mathcal{U}$ , which has  $CEA(\mathbf{c})$  degree, must have c.e. degree?

In particular, as  $\mathcal{U}$  we can take different levels of the Ershov hierarchy.

### Question (Arslanov, Batyrshin, Yamaleev)

Does the construction guarantee that the degree  $CEA(\mathbf{c})$  belongs to the least possible level of the Ershov hierarchy?

### Question (Arslanov, Batyrshin, Yamaleev)

Given low, but non-superlow, c.e. degree  $\mathbf{c}$ . Does there exist  $CEA(\mathbf{c})$  degree which is not of 2-c.e. degree?

## Comments

- What if we try to take all  $\Delta_2^0$ -sets instead of  $\mathcal{U}$ ?
- Then we have to deal with  $\Sigma_2^0$ -sets as well and they ruins the lowness strategies.
- Considering incompleteness strategy we add a freedom (in particular, we can make additional copies of strategies in manner of  $\mathbf{0}'''$ -argument).

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