

An application of computable continuous model theory to a question in proof theory

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www.personal.psu.edu/jmr71/

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Convergence

Convergence theorems

Monotone convergence principle

Let (X, d) be a complete metric space with a linear order $<$ satisfying

$$x < y < z \quad \rightarrow \quad d(x, z) = d(x, y) + d(y, z).$$

Any bounded nondecreasing sequence $(c_n)_{n \in \mathbb{N}}$ converges.

Mean ergodic theorem

Let $(X, \|\cdot\|)$ is a reflexive Banach space with a nonexpansive linear transformation $T: X \rightarrow X$,

$$T(ax + y) = aT(x) + T(y) \quad \text{and} \quad \|T(x)\| \leq \|x\|.$$

Then for $c \in X$, the ergodic averages $\frac{1}{n} \sum_{k < n} T^k(c)$ converge.

Our setup

- Let \mathcal{M} be a complete metric space (X, d) with possible additional structure.
- Let $(c_n)_{n \in \mathbb{N}}$ is a distinguished sequence in X .
- Let P be a property that could hold of $(\mathcal{M}, (c_n)_{n \in \mathbb{N}})$.

Convergence theorem template

If P holds of the pair $(\mathcal{M}, (c_n)_{n \in \mathbb{N}})$, then c_n converges.

Question

For which properties P is the rate of convergence

- **uniform**—exists single rate for all pairs $(\mathcal{M}, (c_n)_{n \in \mathbb{N}})$ satisfying P ?
- **computable**—rate is computable uniformly from $(\mathcal{M}, (c_n)_{n \in \mathbb{N}})$?
- **computably uniform**—exists single computable uniform rate?

Bait-and-switch

- This talk is not about usual Cauchy rates of convergence...

Metastable convergence

Three ways to say converge

The following are all equivalent ways to say that $(c_n)_{n \in \mathbb{N}}$ converges.

- $(c_n)_{n \in \mathbb{N}}$ is Cauchy (**contains a lot of information, but not very uniform**)

$$\underbrace{\forall \varepsilon > 0 \exists m \in \mathbb{N} \forall n, n' \geq m d(c_n, c_{n'}) < \varepsilon.}_{\text{rate of convergence}}$$

- $(c_n)_{n \in \mathbb{N}}$ has finitely many ε -jumps

$$\underbrace{\forall \varepsilon > 0 \exists n \in \mathbb{N} \forall u_0 \leq v_0 \leq u_1 \leq v_1 \leq \dots \leq v_{n-1} \leq v_n}_{\text{rate of convergence}}$$

$$\exists k \in [0, n-1] d(c_{u_k}, c_{v_k}) < \varepsilon.$$

- (Similar to upcrossing bounds and variational bounds.)
- $(c_n)_{n \in \mathbb{N}}$ is metastable (**very uniform, but contains little information**)

$$\underbrace{\forall \varepsilon > 0 \forall F : \mathbb{N} \rightarrow \mathbb{N} \exists m \in \mathbb{N} \forall n, n' \in [m, F(m)] d(c_n, c_{n'}) < \varepsilon.}_{\text{rate of convergence}}$$

Why metastability?

■ Analysis

- Any rate of convergence is better than no rate.
- Rates of metastable convergence are more uniform.
- Better rates may not be known (or even possible?).
- May give simplest or most accessible proof of convergence.
- An alternative to nonstandard analysis.
- Example: Tao's ergodic theorem for multiple commuting averages.

■ Logic

- Metastable rates are computable.
- Metastable convergence theorems are constructive.
- Proof theoretic methods exist to extract metastable bounds: proof mining.
- Closely connected to ultraproducts and nonstandard analysis.
- **Uniform metastable rates can be computed from the statement of the theorem alone! (This talk.)**

Results: prior and new

Let P be a property that could hold of an arbitrary metric structure with a distinguished subsequence: $\mathbb{X} = (X, d, \dots, \{c_n\}_{n \in \mathbb{N}})$. Consider a theorem:

(*) If P holds of \mathbb{X} then c_n converges.

- Kohlenbach. *Some logical metatheorems with applications to functional analysis*. Trans AMS, 2004.
 - If (1) the theorem (*) is provable in $A^\omega[X, d]$,¹ and
 - (2) P is expressible by a \forall -formula (basically a Π_1^0 property), then
 - there exists a computable uniform metastable rate of convergence (uniformly extractable from the proof of (*) in $A^\omega[X, d]$).
- Avigad, Iovino. *Ultraproducts and metastability*. NYJM, 2013.
 - If $\mathcal{C} = \{\mathbb{X} : \mathbb{X} \text{ satisfies } P\}$ is closed under ultraproducts, then
 - there is a uniform rate of metastable convergence.
- R. (This talk)
 - If P is axiomatizable by a set of sentences Σ in continuous logic, then
 - there is a uniform rate of metastable convergence computable from Σ .

¹ $A^\omega[X, d]$ is a type theory extending PA + DC with a type for X and axioms for the metric d .

Main result

Continuous logic

- Continuous logic is a logic for dealing with “metric structures.”
- There have been many variants over the years.
 - Chang and Keisler 1966
 - Henson, et al.
 - etc.
- The current variant is due to Ben Yaacov.
- Ben Yaacov’s version very much resembles first-order logic!
 - Compactness, completeness, Lowenheim-Skolem, ultraproducts, etc.
- See survey article *Model theory for metric structures* for a good introduction.

First-order logic vs. Continuous first-order logic

	First-Order Logic	Continuous First-Order Logic
Universe	set $(M, =)$	comp. bdd. metric space (M, d) ($d(x, y) \leq 1$)
Truth values	T and F	$[0, 1]$ ($0 = \text{true}, 1 = \text{false}$)
Func. symbol	symbol f (arity n)	symbol f (arity n and mod. of cont. $\delta(\varepsilon)$)
Functions	$f^{\mathcal{M}} : M^n \rightarrow M$	$f^{\mathcal{M}} : M^n \rightarrow M$ (obeys mod. of cont. $\delta(\varepsilon)$)
Rel. symbol	symbol R (arity n)	symbol R (arity n and mod. of cont. $\delta(\varepsilon)$)
Relations	$R^{\mathcal{M}} : M^n \rightarrow \{T, F\}$	$R^{\mathcal{M}} : M^n \rightarrow [0, 1]$ (obeys mod. of cont. $\delta(\varepsilon)$)
Connectives	$\odot : \{T, F\}^n \rightarrow \{T, F\}$	$\odot : [0, 1]^n \rightarrow [0, 1]$ (continuous)
Sufficient con.	$\perp, \vee, \wedge, \rightarrow$	$1, \min, \max, \div, x \mapsto x/2$
Quantifiers	$\exists x \varphi(x), \forall x \varphi(x)$	$\min_x \varphi(x), \max_x \varphi(x)$
Formulas	$\varphi(\bar{x})$	$\varphi(\bar{x})$
Statements	Sentence: φ	Conditions: $[\varphi = 0], [\varphi > 0]$
Evaluation	$\varphi^{\mathcal{M}} : M^n \rightarrow \{T, F\}$	$\varphi^{\mathcal{M}} : M^n \rightarrow [0, 1]$
Satisfaction	$\mathcal{M} \models \varphi$ iff $\varphi^{\mathcal{M}} = T$	$\mathcal{M} \models [\varphi = 0]$ iff $\varphi^{\mathcal{M}} = 0$
Axioms/rules	complete, comp. list	complete, computable list
Provability	$\Sigma \vdash \varphi$	$\Sigma \vdash [\varphi > 0]$ (Σ set of conditions $[\psi = 0]$)

Main theorem

Theorem (R.)

Let \mathcal{L} be a computable signature containing constants $\{c_n\}_{n \in \mathbb{N}}$.

Let Σ be a set of \mathcal{L} -conditions (using connectives $1, \min, \max, \div, \cdot/2$).

Assume $(c_n^{\mathcal{M}})_{n \in \mathbb{N}}$ converges for all $\mathcal{M} \models \Sigma$.

Then there is a uniform rate of metastability for $(c_n^{\mathcal{M}})_{n \in \mathbb{N}}$ computable in Σ .

Proof.

- Fix $F: \mathbb{N} \rightarrow \mathbb{N}$ and $\varepsilon = 2^{-k}$, our goal is to find ℓ such that

$$\exists m < \ell \forall n, n' \in [m, F(m)] d^{\mathcal{M}}(c_n^{\mathcal{M}}, c_{n'}^{\mathcal{M}}) < \varepsilon.$$

- The quantifiers are bounded, so this can be written as the \mathcal{L} -condition

$$\underbrace{\max_{m < \ell} \min_{n, n' \in [m, F(m)]} \left(1/2^k \div d(c_n, c_{n'}) \right)}_{\varphi_\ell} > 0.$$

Main theorem (cont.)

Proof.

- Fix $F: \mathbb{N} \rightarrow \mathbb{N}$ and $\varepsilon = 2^{-k}$, our goal is to find ℓ such that for all $\mathcal{M} \models \Sigma$,

$$\exists m \leq \ell \forall n, n' \in [m, F(m)] d^{\mathcal{M}}(c_n^{\mathcal{M}}, c_{n'}^{\mathcal{M}}) < \varepsilon.$$

- The quantifiers are bounded, so this can be written as the \mathcal{L} -condition

$$\underbrace{\max_{m < \ell} \min_{n, n' \in [m, F(m)]} \left(1/2^k \div d(c_n, c_{n'}) \right)}_{\varphi_\ell} > 0.$$

- Notice φ_ℓ is computable in F, k , and ℓ .
- For all $\mathcal{M} \models \Sigma$, since $(c_n^{\mathcal{M}})_{n \in \mathbb{N}}$ converges, there exists ℓ s.t. $\mathcal{M} \models [\varphi_\ell > 0]$.
- By the compactness theorem, there is some ℓ such $\Sigma \models [\varphi_\ell > 0]$.
- By the completeness theorem, $\Sigma \vdash [\varphi_\ell > 0]$ for some ℓ .
- Compute ℓ by searching for a finite proof of $\Sigma \vdash [\varphi_\ell > 0]$ for some ℓ . □

Closing Thoughts

Summary

- There are many compatible logical tools for investigating theorems in analysis:
 - type theory and the dialectic interpretation,
 - ultraproducts, and
 - continuous logic.
- Usual (computability theoretic, proof theoretic, model theoretic) methods from first order logic extend nicely to continuous logic.
- There is a lot of potential to investigate computable continuous model theory; it is a nice merger of computable analysis and computable model theory.
- Can continuous logic be applied to proof mining in a useful way?
- Can logic methods be used to study other types of rates of convergence?
- **See my NERDS talk in two weeks for more details and examples.**

Thank You!

These slides will be available on my webpage:

<http://www.personal.psu.edu/jmr71/>

Or just Google™ me, “Jason Rute”.

P.S. I am on the job market.