

An Unusual Reflection Principle for Self Justifying Logics

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1. Overview

Gödel's 2nd Incomplet. Theorem indicates strong formalisms cannot verify their own consistency

But Humans Intuitively Appreciate their Own Consistency

Topic of our 64 Page Paper: What kinds of systems are Adequately Weak to formalize **some type (?)** of knowledge of their own consistency?

Research in New Technical Report and Six Prior Articles in JSL and APAL Has Sought to:

- 1 Develop New Generalizations of Second Inc Theorem
- 2 Formalize **Unusual** "Boundary-Case Exceptions" to It.
- 3 Produce Tightest Possible Match Between Items 1 + 2.

2. Background Literature (summarized in 3 slides)

Definition: Axiom System β called **Self Justifying** relative to Deduction Method d when :

- 1 one of β 's formal theorems states d 's deduction method, applied to axiom system β , is consistent.
- 2 and the axiom system β is **also actually** consistent.

$\forall \alpha \forall d$ Kleene (1938), Rogers (1966) & Jeroslow (1971) noted **Easy To Construct** axiom system $\alpha^d \supseteq \alpha$ satisfying **Requirement 1** i.e. set $\alpha^d = \alpha \cup \text{SelfCons}(\alpha, d)$ (defined below)

“There is no proof (using d 's deduction method) of $0 = 1$ from the **Union** of system α with **this** sentence (**looking at itself**)”

Above Well Defined But Catch is α^d **Usually Fails** Item 2.

- i.e. α^d is inconsistent via a Gödel diagonalization paradigm.
- Thus prior to Willard (1993), this topic mostly shunned.

3. More Background Literature

Definition: Let α denote axiom system lacking Induction Principle
Then $\Psi(x)$ called α -Initial Segment iff α can prove:

$$\Psi(0) \quad \text{AND} \quad \forall x \quad \Psi(x) \rightarrow \Psi(x + 1) \quad (1)$$

Pudlák 1985: All axiom systems of finite cardinality have Initial Segments Ψ where α can verify its Herbrand and Semantic Tableaux Consistency for every x satisfying $\Psi(x)$

- **Intuition:** All integers x satisfy $\Psi(x)$ BUT α NOT KNOW THIS !
- Above Result does not generalize for Hilbert Deduction

Kreisel-Takeuti (1974) Earliest Local-Consistency Result:

- Showed Second-Order Generalization of Cut-Free Deduction Can Verify Its Own Consistency.
- Sets Ψ (in Equation 1) = Dedekind's Definition of Integers

Verbrugge-Visser (1994) developed analogous arithmetic reflection principles using local consistency constructs.

- Visser (2005) discusses this topic further and summarizes Harvey Friedman's Ohio State 1979 Tech Report

4. Generalizations of Second Inc Theorem

- **Bezboruah-Shepherdson 1976**: Showed some Gödel encodings of Robinson's Q CANNOT VERIFY their Hilbert consistency.
- **Pudlák 1985**: Generalized Above for all Gödel encodings of proofs and for **All Initial Segments** (defined on prior slide) when Hilbert Deduction Present.
- **Wilkie-Paris 1987** : showed $I\Sigma_0 + \text{Exp}$ CANNOT PROVE Hilbert Consistency of Q,
- **Solovay (1994 Private Com.)** : Showed **NO SYSTEM** (weaker than Q) Recognizing **MERELY SUCCESSOR** as total function can VERIFY its Hilbert Consistency.
- **W— 2002-2009** : generalized work of Adamowicz-Zbierski to show **THREE DIFFERENT ENCODINGS** of $I\Sigma_0$ CANNOT PROVE their semantic tableaux consistency.

Hence Self-Justifying Formalisms **Always Contain** weaknesses.

5. Main Perspective of Willard's 1993-2009 Research

Notation: $Add(x, y, z)$ and $Mult(x, y, z)$ are 3-way atomic predicates employed by our axiom systems.

Definitions: An axiom system α is

- **Type-A** iff it contains Equation 1 as axiom:
- **Type-M** iff it contains $1 + 2$ as axiom:
- **Type-S** iff it can prove (3) **BUT NOT PROVE** (1) NOR (2) :

$$\forall x \forall y \exists z \quad Add(x, y, z) \quad (1)$$

$$\forall x \forall y \exists z \quad Mult(x, y, z) \quad (2)$$

$$\forall x \exists z \quad Add(x, 1, z) \quad (3)$$

Combined Result of Pudlak, Solovay, Nelson, Wilkie-Paris:

- No natural Type-S system can recognize its Hilbert consistency:

Our Main Prior Results about this Subject:

- 1 Some Type-A prove all PA's π_1 theorems and their semantic tableaux consistency
- 2 Most Type-M axiom systems **UNABLE to JUSTIFY** their semantic tableaux consistency.

6. Limitations Upon Self Justifying Systems

- 1 Pudlak (1985) + Solovay (1994) (combined with Nelson + Wilkie-Paris) implies self-justification collapses when Hilbert Deduction is present for most systems recognizing Successor as total function.
- 2 JSL(2002)+ APAL(2007) indicates Semantic Tableaux Self Justification collapses when Multiplication recognized as Total Function.
- 3 FOL-2004 Paper showed that while JSL 2005 could add a π_1 and Σ_1 modus ponens rule to our semantic tableaux evasions of Second Incompleteness Theorem, **Same NOT TRUE with π_2 and Σ_2 modus ponens rules.**

Next Three Slides Have **GOOD NEWS** despite Items 1-3:

Self-Justifying Systems Support Unusually Robust Reflection Principles.

Thus Bad News from Items 1-3 **Not Fully Dismal !**

7. New Perspective about Reflection Principles

Def: $\text{Reflect}_{\alpha,D}(\Psi)$ denotes sentence Ψ 's reflection principle under the axiom system α and deduction method D i.e.

$$\forall p \{ \text{Prf}_{\alpha,D}(\ulcorner \Psi \urcorner, p) \Rightarrow \Psi \} \quad (4)$$

Löb's Theorem: If $\alpha \supset \text{Peano Arith}$ then α cannot prove $\text{Reflect}_{\alpha,D}(\Psi)$ except in trivial case where it can prove Ψ .

Gödel's Anti-Reflection Theorem: No reasonable axiom system α can prove $\text{Reflect}_{\alpha,D}(\Psi)$ for all π_1 sentences.

i.e. Difficulties always arise because Gödel Sentences declaring "There is no proof of me" have π_1 encodings.

Surprising Fact: Self-Justifying Systems Support "Transformed" π_1 Reflection Principles Despite Above 2 Theorems, i.e.

$$\forall p \{ \text{Prf}_{\alpha,D}(\ulcorner \Psi \urcorner, p) \Rightarrow \Psi^T \} \quad (5)$$

where T is isomorphism mapping π_1 sentences into π_1 sentences such that $\Psi \leftrightarrow \Psi^T$ holds in Standard Model.

8. Two New Theorems About Reflection Principles

Def: Ax System α is **Level(1^D) Consistent** iff α **UNABLE TO PROVE** under deduction method D **BOTH** some π_1 sentence and its negation.

Theorem 6.12 If α can formally verify its own Level(1^D) Consistency
Then there exists some T where α can verify (6)'s "Transformational"
Reflection Principle for **All** π_1 sentences Ψ **simultaneously**.

$$\forall p \{ \text{Prf}_{\alpha,D}(\ulcorner \Psi \urcorner, p) \Rightarrow \Psi^T \} \quad (6)$$

Intuition Behind Theorem 6.12 : The identity $\Psi \leftrightarrow \Psi^T$ holds in Standard Model, **BUT** α **UNABLE** to verify it.

Theorem E.1 If Ax System α unable to prove its own consistency (i.e. satisfies Second Inc.Theorem) then α **UNABLE TO VERIFY** (6)'s Transform Reflection Principle for **All** π_1 sentences Ψ **simultaneously**.

Proof Sketch: All conventional axiom systems can refute all false π_1 sentences. Hence if Ψ false then α can refute both Ψ and Ψ^T . But then α could use (6)'s reflection principle to confirm its own consistency. **Latter impossible** because contradicts Theorem 6.12's hypothesis. \square

9. Mysterious Two Sentences in Gödel's 1931 Paper

Most Surprising Two Sentences in Gödel's Paper:

- *"It must be expressly noted that Theorem XI (i.e the Second Inc Theorem) represents no contradiction of the formalistic standpoint of Hilbert. For this standpoint presupposes only the existence of a consistency proof by finite means, and there might conceivably be finite proofs which cannot be stated in ... "*

Our Interpretation of Gödel's Statement • :

- 1 We agree with most logicians that Gödel was excessively cautious in Statement • because history has proven the Second Inc Theorem to be a 95 % Robust Result from a "Consistency Perspective".
- 2 However, Gödel's Statement • is **QUITE SIGNIFICANT** from a "Reflection Perspective" because π_1 Transform Reflection explains how Thinking Beings acquire motivation to cogitate.

$$\forall p \{ \text{Prf}_{\alpha,D}(\ulcorner \Psi \urcorner, p) \Rightarrow \Psi^T \} \quad (7)$$

10. Concluding Remarks

Wide Significance of Gödel's 2nd Incomp Theorem illustrated by:

- Its generalization using 1939 Hilbert-Bernays Derivation Conditions
- Solovay's 1994 Extension of Pudlák's 1985 Work:

No Axiom System *viewing successor as a total function can justify its own Hilbert consistency.*

Above Precludes **many but not all** uses of **"I am consistent"** axioms:

- 1 This is because Reflection Principles explain how Thinking Beings Motivate Themselves to Cogitate
- 2 This use of Reflection Principles **Is Very Helpful**, EVEN IF it does not formalize a **STRONG RESPECT** where systems confirm their own consistency.

Many Other Results at <http://arxiv.org/abs/1108.6330>.

- Purpose of this Talk was to be pointer to 64-page report
- Latter Both Unifies and Extends our Prior Results