

# Seven characterizations of the cototal enumeration degrees

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## Motivation from symbolic dynamics

### Definition

- A *subshift* is a nonempty closed set  $X \subseteq 2^\omega$  such that if  $a\alpha \in X$  then  $\alpha \in X$ .
- $X$  is *minimal* if there is no  $Y \subset X$ , such that  $Y$  is a subshift.

Given a minimal subshift  $X$ , we would like to characterize the set of Turing degrees that compute members of  $X$ .

### Definition

The *language* of subshift  $X$  is the set

$$L_X = \{\sigma \in 2^{<\omega} \mid \exists \alpha \in X (\sigma \text{ is a subword of } \alpha)\}.$$

- 1 If  $X$  is minimal and  $\sigma \in L_X$  then for every  $\alpha \in X$ ,  $\sigma$  is a subword of  $\alpha$ . So every element of  $X$  can enumerate the set  $L_X$ .
- 2 If we can enumerate  $L_X$  then we can compute a member of  $X$ .

## The enumeration degrees and cototal sets

### Definition

$A \leq_e B$  if every enumeration of  $B$  can compute an enumeration of  $A$ .

The *enumeration-cone* of a set  $A$  is the set of Turing degrees that can enumerate  $A$ .

The enumeration-cone of  $L_X$  is the set of Turing degrees that compute members of  $X$ .

(Jaendel:) If we can enumerate the set of *forbidden words*  $\overline{L_X}$  then we can enumerate  $L_X$ .

So  $L_X \leq_e \overline{L_X}$ .

### Definition

A set  $A$  is *cototal* if  $A \leq_e \overline{A}$ . An enumeration degree is cototal if it contains a cototal set.

## Examples of cototal enumeration degrees

### Proposition

Every total e-degree is cototal.

$$A \oplus \bar{A} \equiv_1 \overline{A \oplus \bar{A}} = \bar{A} \oplus A.$$

### Proposition

Every  $\Sigma_2^0$  e-degree is cototal.

Let  $A$  be  $\Sigma_2^0$ . Consider the set  $K_A = \bigoplus_{e < \omega} \Gamma_e(A)$ . Then  $A \equiv_e K_A$  and

$$\bar{K}_A = \bigoplus_{e < \omega} \overline{\Gamma_e(A)} \geq_e \bar{K} \geq_e A \equiv_e K_A.$$

## Characterization I: The skip operator

Note, that  $A \leq_e B$  if and only if  $\overline{K_A} \leq_1 \overline{K_B}$ .

### Definition (AGKLMSS)

The *skip* of  $A$  is the set  $A^\diamond = \overline{K_A}$ . The *skip* of a degree is  $\mathbf{d}_e(A)^\diamond = \mathbf{d}_e(A^\diamond)$ .

Recall, that the *enumeration jump* of  $A$  is defined by  $A' = K_A \oplus \overline{K_A}$ . So for every enumeration degree  $\mathbf{a}$  we have that  $\mathbf{a}' = \mathbf{a} \vee \mathbf{a}^\diamond$ .

### Theorem (AGKLMSS)

Let  $S \geq_e \emptyset'$ . There is a set  $A$  such that  $A^\diamond \equiv_e S$ .

### Proposition (AGKLMSS)

A degree  $\mathbf{a}$  is cototal if and only if  $\mathbf{a} \leq \mathbf{a}^\diamond$  (if and only if  $\mathbf{a}^\diamond = \mathbf{a}'$ ).

## Characterization II: A topological perspective

Miller introduced the *continuous degrees*  $\mathcal{D}_r$  to compare the complexity of points in computable metric spaces.

The continuous degrees embed into  $\mathcal{D}_e$ . In fact,  $\mathcal{D}_T \subset \mathcal{D}_r \subset \mathcal{D}_e$ .

### Proposition (AGKLMSS)

The every continuous enumeration degree is cototal.

Kihara and Pauly extend Miller's idea to points in arbitrary *represented* topological spaces. They define the *point degree spectrum* of a represented space.

- 1  $Spec(\{0, 1\}^\omega) = Spec(\omega^\omega) = \mathcal{D}_T$ ;
- 2  $Spec([0, 1]^\omega) = Spec(C([0, 1])) = \mathcal{D}_r$ ;
- 3  $Spec(S^\omega) = \mathcal{D}_e$ , where  $S$  is the Sierpinski space.

### Theorem (Kihara)

The cototal e-degrees are the elements of the point degree spectra of all sufficiently effective second countable  $G_\delta$  spaces. (Every closed set is  $G_\delta$ ).

## Characterization III and IV: Graph theory

### Definition (Carl von Jaenisch)

Let  $G = (V, E)$  be a graph. A set  $M \subseteq V$  is *independent*, if no two members of  $M$  are edge related.  $M$  is *maximal* set, if every  $v \in V$  is either in  $M$  or edge related to a vertex in  $M$ .

### Theorem (AGKLMSS)

An enumeration degree is cototal if and only if it contains the complement of a maximal independent set for the graph  $\omega^{<\omega}$ .

### Theorem (McCarthy)

An enumeration degree is cototal if and only if it contains the complement of a maximal antichain in  $\omega^{<\omega}$ .

## Characterization V: E-pointed trees

### Definition (Montalbán)

A tree  $T \subseteq 2^{<\omega}$  is *enumeration pointed* if it has no dead ends and every infinite path  $f \in [T]$  enumerates  $T$ .

### Theorem (Montalbán)

A degree spectrum is never the Turing-upward closure of an  $F_\sigma$  set of reals in  $\omega^\omega$ , unless it is an enumeration-cone.

### Theorem (McCarthy)

An enumeration degree is cotal if and only if it contains a (uniformly) enumeration pointed tree.

### Corollary

A degree spectrum is the Turing-upward closure of an  $F_\sigma$  set of reals in  $\omega^\omega$  if and only if it is the enumeration-cone of a cotal e-degree.



## Characterization VI: Minimal subshifts

The enumeration degree of the language  $L_X$  of a minimal subshift  $X$  characterizes the set of Turing degrees of members of  $X$ .

(Jaendel:)  $L_X$  is a cototal set.

### Theorem (McCarthy)

Every cototal enumeration degree is the degree of the language of a minimal subshift.

## Characterization VII: Good enumeration degrees

### Definition (Lachlan, Shore)

A uniformly computable sequence of finite sets  $\{A_s\}_{s < \omega}$  is a *good approximation* to a set  $A$  if:

$$\mathbf{G1} \quad (\forall n)(\exists s)(A \upharpoonright n \subseteq A_s \subseteq A)$$

$$\mathbf{G2} \quad (\forall n)(\exists s)(\forall t > s)(A_t \subseteq A \Rightarrow A \upharpoonright n \subseteq A_t).$$

An enumeration degree is *good* if it contains a set with a good approximation.

- 1 Good e-degrees cannot be tops of empty intervals.
- 2 Total enumeration degrees and enumeration degrees of  $n$ -c.e.a. sets are good.

## Characterization VII: Good enumeration degrees

### Theorem (Harris; Miller, S)

The good enumeration degrees are exactly the cototal enumeration degrees.

If  $A$  has a good approximation then

$$A \leq_e \{ \langle x, s \rangle \mid (\forall t > s)(A_t \subseteq A \Rightarrow x \in A) \} \leq_e A^\diamond.$$

Every uniformly enumeration pointed tree has a good approximation.

### Theorem (Miller, S)

The cototal enumeration degrees are dense.

If  $V <_e U$  are cototal and  $U$  has a good approximation we can build  $\Theta$  such that  $\Theta(U)$  is the complement of a maximal independent set and

$$V <_e \Theta(U) \oplus V <_e U.$$

The end

Thank you!