## Qualifying Exam

#### LOGIC

#### August 26, 1982

Do FOUR of the following problems, at most two elementary.

### I. Elementary Questions

- 1. Classify each of the following classes (with regard to first-order logic) as
  - a) finitely axiomatizable (give the axioms)
  - b) axiomatizable but not finitely axiomatizable (give the axioms and a proof that no finite set of axioms suffices).
  - c) not axiomatizable (give proof).
  - i) The class of finite linearly ordered sets, (A,<)
  - ii) The class of infinite linearly ordered sets, (A,<)
  - iii) The class of densely linearly ordered sets, (A,<)
  - iv) The class of well-ordered sets, < A , <> .
- 2. Prove the following or give a counter-example. Let S be a decidable (i.e., recursive) set of sentences in propositional logic, using proposition letters  $P_n$  (n  $\epsilon$   $\omega$ ). Then  $\{\varphi\colon S \vdash \varphi\}$  is decidable.
- 3. Give a counterexample to the following statement:

If  $\alpha, \gamma$  are ordinals and  $A \subset \gamma$ ,  $A \cong \alpha$  and  $(\gamma-A) \cong \alpha$ , then  $\gamma = \alpha + \alpha$  or  $\gamma = \alpha$ .

II. Model Theory.

- Which of the following properties of complete theories and/or models are preserved under the operation of taking reducts (explain):
  - i) % -categoricity;
- iv) being prime;
- ii) <sup>ℵ</sup>₁-categoricity;
- v) being saturated;

iii)  $\omega$ -stability;

- vi) having only finitely many countable models.
- 2. Assume T is a complete theory,  $\{\Gamma_{\eta} | \eta \epsilon 2^{<\omega} \}$  and  $\{A_{\eta} | \eta \epsilon 2^{<\omega} \}$  are types and models of T respectively, satisfying:
  - i) if  $\eta \subset \xi$  , then  $A_\xi$  realizes  $\Gamma_\eta$  and  $\Gamma_\eta \subset \Gamma_\xi$  ; and

Prove then that T has 2 pairwise non-isomorphic countable models.

- 3. Suppose M < M' , a,b  $\in |M|$  , c  $\in |M'|$  , p,q types over M ,  $\varphi(\mathbf{x},\mathbf{a}) \text{ a formula with parameters from } |M| \cup \{\mathbf{a}\} \text{ satisfy:}$ 
  - i) a realizes p and b realizes q in M';
  - ii)  $p(\bar{x}) \cup q(\bar{y})$  is a complete type; and
  - iii)  $\varphi(\mathbf{x},\mathbf{a})$  is a complete formula for the type over  $|\mathbf{M}| \cup \{\mathbf{a}\}$  realized by c .

Prove then that  $\varphi(\mathbf{x}, \mathbf{a})$  is also a complete formula for the type that  $\mathbf{c}$  realizes over  $|\mathbf{M}| \cup \{\mathbf{a}, \mathbf{b}\}$ . [Hint: Otherwise there is  $\mathbf{c}' \in \mathbf{M}'' \geq \mathbf{M}'$  realizing the same type as  $\mathbf{c}$  over  $|\mathbf{M}| \cup \{\mathbf{a}\}$ , but not over  $|\mathbf{M}| \cup \{\mathbf{a}, \mathbf{b}\}$ . Proceed from there!].

# III. Recursion Theory.

- 1. Prove that there is no recursive f satisfying:
  - i)  $W_{f(n)}$  codes a well order,  $n < \omega$ ;
  - ii) if  $W_n$  is a well order, then  $W_n = W_{f(n)}$ ,  $n < \omega$ .  $[A \subseteq \omega \text{ codes a well order if } \langle \operatorname{ran} A \cup \operatorname{dom} A, \{(x,y) | \langle x,y \rangle \in A \} \rangle$  is a well order, where  $\operatorname{ran} A = \{y | \exists x (\langle x,y \rangle \in A) \}$  and  $\operatorname{dom} A = \{x | \exists y (\langle x,y \rangle \in A) \} \}.$

[Hint: One way to do this uses the recursion theorem].

- 2. Suppose C is a non-recursive  $\Delta_2^0$  set. Prove that there is a simple set A such that  $C \not\leftarrow_T A$ . Hint: Use the limit lemma to represent C as  $\lim_{s \to \infty} C_s$ , where  $\{C_s\}_{s < \omega}$  is a recursive sequence.
- 3. Suppose A is a countable, saturated structure realizing recursive types  $\Gamma_1, \Gamma_2 \ . \quad \text{Prove that there is a decidable structure realizing} \quad \Gamma_1 \quad \text{and} \quad \Gamma_2 \ .$
- 4. Suppose  $\{A_s\}_{s<\omega}$  is an effective enumeration of A , where each  $A_s$  is recursive, and  $\forall e\{[\exists^\infty s \{e\}^N_s(e)\downarrow\} \rightarrow \{e\}^A(e)\downarrow\}$ .

  Prove A is low, i.e. A' = 0'.

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# IV. Set Theory.

- 1. Let (A,<) be a linear order such that  $\forall x \in A$   $(x \cong A \text{ or } A-x \cong A)$ . Show A is a well order or an inverse well order.
- 2. Call F a  $\underline{\mathrm{mad}\ f}$  (maximal almost disjoint family of subsets of  $\omega_1$ ) iff  $i) \quad F \subseteq P(\omega_1) \ .$ 
  - ii)  $\forall x \in F \ (|x| = \omega_1)$ .
  - iii)  $\forall x, y \in F \quad (x \neq y \rightarrow |x \cap y| < \omega_1)$ .
  - iv) F is maximal with respect to (i) (iii) .

Assume: M is a countable transitive model of ZFC ,  $F \in M$  , and  $(F \text{ is a mad f})^M$  . Show that F remains a mad f in any ccc forcing extension of M .

- 3. If  $f,g \in \omega^{\omega}$ , we say f < g iff  $\{n: g(n) \leq f(n)\}$  is finite. Let  $\alpha$  be the least ordinal which is not the type of a chain in the partial order,  $(\omega^{\omega}, < g^*)$ . Prove that  $\alpha$  is a regular cardinal.
- 4. Assume that there is an ordinal  $\alpha$  such that  $R(\alpha) \models ZFC$ , and let  $\alpha$  be the smallest such. Prove that  $\alpha$  is a strong limit cardinal and that  $cf(\alpha) = \omega$ .