

Toward deciding the $\forall\exists$ -theory of the Σ_2^0 -enumeration degrees

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(joint work with Goh, Ng and M. Soskova)

Most “natural” degree structures \mathcal{D} are very complicated partial orders and usually follow this pattern:

- The first-order theory of the partial order \mathcal{D} is undecidable. In fact, it is usually as complicated as second-order arithmetic (for global degree structures) or first-order arithmetic (for countable local degree structures).

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Definition

$A \leq_e B$ if there is an enumeration operator Φ with $A = \Phi(B)$, i.e., there is a c.e. set Φ of pairs (x, F) (of numbers x and finite sets F) denoting that for all x , $x \in A$ iff there is $(x, F) \in \Phi$ with $F \subseteq B$.

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In particular, we will focus on the degree structure \mathcal{S}_e of the enumeration degrees of the Σ_2^0 -sets, which coincides with the enumeration degrees $\mathbf{a} \leq \mathbf{0}'_e$. They form a densely ordered countable upper semilattice with least element $\mathbf{0}_e$ (the degree of the c.e. sets) and greatest element $\mathbf{0}'_e$ (the degree of \overline{K}).

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For \mathcal{S}_e , the \exists -theory is decidable by Lagemann (1972), whereas the $\exists\forall\exists$ -theory is undecidable by Kent (2006).

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However, the decidability of the $\forall\exists$ -theory of \mathcal{S}_e remains open.

Deciding the $\forall\exists$ -theory of a degree structure \mathcal{D} amounts to giving a uniform decision procedure to the following

Algebraic Problem (for deciding the $\forall\exists$ -theory of \mathcal{D})

Given finite partial orders \mathcal{P} and $\mathcal{Q}_i \supseteq \mathcal{P}$ (for $i \leq n$), does every embedding of \mathcal{P} into \mathcal{D} extend to an embedding of \mathcal{Q}_i into \mathcal{D} for some $i \leq n$ (where i may depend on the embedding of \mathcal{P})?

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Two major subproblems of the $\forall\exists$ -theory of \mathcal{S}_e are decidable:

Extension of Embeddings Problem

Given finite partial orders \mathcal{P} and $\mathcal{Q} \supseteq \mathcal{P}$, does every embedding of \mathcal{P} into \mathcal{S}_e extend to an embedding of \mathcal{Q} into \mathcal{S}_e ?
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Lattice Embeddings Problem

Which finite lattices can be embedded into \mathcal{S}_e (preserving join and meet)? (Lempp/Sorbi 2002: all finite lattices embed)

The main technical obstacles to deciding the $\forall\exists$ -theory of \mathcal{S}_e showed up first in the following

Theorem (Ahmad 1989 (cf. Ahmad/Lachlan 1998))

- 1 There is an *Ahmad pair* of Σ_2^0 -enumeration degrees (\mathbf{a}, \mathbf{b}) , i.e., there are incomparable degrees \mathbf{a} and \mathbf{b} such that any degree $\mathbf{v} < \mathbf{a}$ is $\leq \mathbf{b}$.

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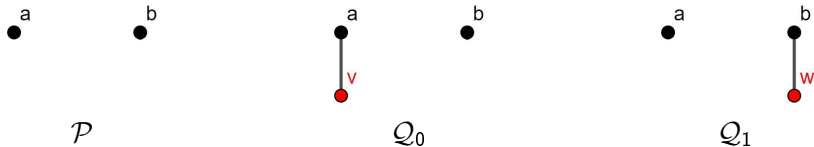
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These are examples of $\forall\exists$ -statements blocking $\mathcal{P} \subset \mathcal{Q}_0$ but not both $\mathcal{P} \subset \mathcal{Q}_0$ and $\mathcal{P} \subset \mathcal{Q}_1$:



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Technical Questions

- 1 Is there an *Ahmad triple* of Σ_2^0 -enumeration degrees, i.e., are there degrees \mathbf{a} , \mathbf{b} and \mathbf{c} such that (\mathbf{a}, \mathbf{b}) and (\mathbf{b}, \mathbf{c}) form Ahmad pairs?

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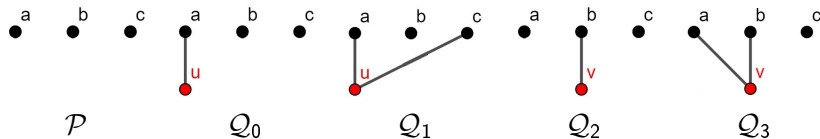
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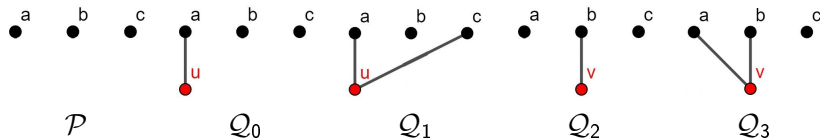


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For many years, I believed the answers to both to be “yes”.

However, the answer to both questions is “no”:

Theorem (Goh, Lempp, Ng, M. Soskova, to appear)

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- 2 But there is a *weak Ahmad triple*, i.e., there are pairwise incomparable Σ_2^0 -enumeration degrees \mathbf{a} , \mathbf{b} and \mathbf{c} such that (\mathbf{a}, \mathbf{b}) and (\mathbf{a}, \mathbf{c}) do not form Ahmad pairs but any degree $\mathbf{v} < \mathbf{a}$ is $\leq \mathbf{b}$ or $\leq \mathbf{c}$.

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Theorem (Kalimullin, Lempp, Ng, Yamaleev, submitted)

There is no cupping Ahmad pair.

The proof turns out to be a non-uniform finite-injury(!) argument.

Given the difficulty of the overall problem of deciding the $\forall\exists$ -theory, we are currently concentrating on the following subproblem:

1-Point Extensions of Antichains

Decide, given a finite antichain $\mathcal{P} = \{a_0, \dots, a_n\}$ and 1-point extensions $\mathcal{Q}_S = \{a_0, \dots, a_n, x_S\}$ and $\mathcal{Q}^T = \{a_0, \dots, a_n, x^T\}$ for some nonempty subsets $S, T \subseteq \{0, \dots, n\}$ (where $x_S < a_i$ iff $i \in S$; and $x^T > a_i$ iff $i \in T$),

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We have now found a (complicated) complete characterization for the above subproblem involving only extensions Q_S .

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We have no working conjecture that combines the \mathcal{Q}_S and the \mathcal{Q}^T .

Main Theorem (Goh, Lempp, Ng, M. Soskova, in preparation)

Fix $n > 0$ and $\mathcal{S} \subseteq \mathcal{P}(\{0, \dots, n\}) - \{\emptyset\}$.

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Then some embedding of \mathcal{P} into \mathcal{S}_e cannot be extended to an embedding of \mathcal{Q}_S for any $S \in \mathcal{S}$ (" S can be blocked") iff (*) holds:

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$\{0, \dots, n\} \neq \bigcup S$: Fix $j \in \{0, \dots, n\} - \bigcup S$ and make each \mathbf{a}_k (for $k \neq j$) form an Ahmad pair with \mathbf{a}_j .

The most difficult condition of (*) concerns the assignment $\nu : S_0 \rightarrow \mathcal{P}(S_1) - \{\emptyset\}$ satisfying

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Harder example showing the second bullet is needed:

\mathbf{a}_0 and \mathbf{a}_1 both form an Ahmad pair with \mathbf{a}_2 , and \mathbf{a}_0 and \mathbf{a}_1 form a minimal pair; so $S_0 = \{0, 1\}$ and $\nu : 0, 1 \mapsto \{2\}$, namely, $\mathcal{S} = \{\{0\}, \{1\}, \{0, 1, 2\}\}$ and even $\mathcal{S} = \{\{0\}, \{1\}, \{0, 1\}, \{0, 1, 2\}\}$ can be blocked.

The most difficult condition of (*) concerns the assignment $\nu : S_0 \rightarrow \mathcal{P}(S_1) - \{\emptyset\}$ satisfying

- for each $i \in S_0$, $\{i\} \cup \nu(i) \notin \mathcal{S}$, and
- for each $F \subseteq S_0$ with $|F| > 1$, we have $\bigcap \{\nu(i) \mid i \in F\} \notin \mathcal{S}$:

Easy example showing the first bullet is needed:

\mathbf{a}_0 forms an Ahmad pair with \mathbf{a}_1 ;

so $S_0 = \{0\}$ and $\nu : 0 \mapsto \{1\}$, namely, $\mathcal{S} = \{\{0\}\}$ can be blocked, but $\{\{0\}, \{0, 1\}\}$ cannot.

Harder example showing the second bullet is needed:

\mathbf{a}_0 and \mathbf{a}_1 both form an Ahmad pair with \mathbf{a}_2 , and \mathbf{a}_0 and \mathbf{a}_1 form a minimal pair; so $S_0 = \{0, 1\}$ and $\nu : 0, 1 \mapsto \{2\}$, namely, $\mathcal{S} = \{\{0\}, \{1\}, \{0, 1, 2\}\}$ and even $\mathcal{S} = \{\{0\}, \{1\}, \{0, 1\}, \{0, 1, 2\}\}$ can be blocked.

But: Note that the first bullet fails for $\mathcal{S} = \{\{0\}, \{1\}, \{0, 2\}\}$, so this cannot be blocked.

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Suppose $S_0 \neq \emptyset$ and $\{0, \dots, n\} = \bigcup \mathcal{S}$.

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Suppose \mathbf{a} , \mathbf{b}_i and $\mathbf{c}_{i,j}$ (for $i < m$ and $j < n_i$) are degrees with $\mathbf{a} \not\leq \mathbf{b}_i$ and $\mathbf{b}_i \not\leq \mathbf{c}_{i,j}$ for all i and j .

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Now suppose \mathcal{P} embeds via degrees \mathbf{a}_i . For each $i \in S_0$, fix nonzero $\mathbf{v}_i < \mathbf{a}_i$ with $\mathbf{v}_i \not\leq \mathbf{a}_k$ for all $k \in S_0 - \{i\}$, and set $\nu(i) = \{j \in S_1 \mid \mathbf{v}_i \leq \mathbf{a}_j\}$, so $\{i\} \cup \nu(i) \notin \mathcal{S}$ (and $\nu(i) \neq \emptyset$).

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On the other hand, for $F \subseteq S_0$ with $|F| > 1$, set $\mathbf{v}_F = \bigcup_{i \in F} \mathbf{v}_i$, and so $\mathbf{v}_F < \mathbf{a}_j$ iff $j \in \bigcap \{\nu(i) \mid i \in F\}$ (and $\bigcap \{\nu(i) \mid i \in F\} \notin \mathcal{S}$).

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So ν is an assignment as desired.

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$$\mathcal{I}_j : A_j \neq W \quad (\text{if } j \in S_1 - \bigcup_{i \in S_0} \nu(i))$$

Thanks!