

# Low levels of the arithmetical hierarchy and computable reductions on $\omega$

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Surveying results of many people.

A lot of work has been focused on the structure of ceers, including:

- There is a universal degree, which appears naturally: Provable equivalence in PA, isomorphism of finite presentations of groups, word problems of some groups, equivalence relations where the classes are uniformly effectively inseparable.
- Ceers with finitely many classes form an initial segment  $\mathcal{I}$ .
- There are ceers which are not above  $=^\omega$  (usually called Id). We call these *dark*. This is a failure of the analog of Silver's theorem.
- There are infinitely many ceers which are minimal over  $\mathcal{I}$ .
- We have some descriptions of when pairs of ceers have (or don't have) a join or a meet.
- Every degree has a strong minimal cover (some only 1, some countably many)

- $\omega^{<\omega}$  embeds as an initial segment of the degrees (sending the empty string to Id).
- The degree structure of Ceers interprets  $(\mathbb{N}, +, \cdot)$  and so the theory is as complicated as possible. Also, the degree structure of the Light ceers, also the degree structure of the Dark ceers. Also, each of these  $\not\leq \mathcal{I}$ .
- The collection of 1-dimensional ceers  $R_X$  for  $X \subseteq \omega$  embeds the 1-degrees of (infinite) c.e. sets.

## Definition (The Halting Jump operator on ceers)

Given a ceer  $X$ , define  $X'$  by  $i \in X' \iff j \in X \text{ and } \phi_i(i) \downarrow \text{ and } X \leq \phi_j(j) \downarrow$ .

- $X' \geq X$  for all  $X$ .
- $X' > Y'$  iff  $X > Y$ .
- $X' \equiv X$  if and only if  $X$  is universal.
- $X' \leq A \oplus B$  implies  $X' \leq A$  or  $X' \leq B$ .

- There is a universal co-ceer  $\pi$ .
- The only ceer which is below a co-ceer is Id, and the ones with finitely many classes.
- Every co-ceer is light (i.e. above Id).

Everything about ceers relativizes (some care needed: Relativizations include  $0'$ -reductions).

- There is universal  $\Sigma_2^0$ -equivalence relation.
- There are dark ones.
- There are the 1-dimensional ones (closed downwards)

We haven't really considered what the halting jump looks like here. e.g., What are there other fixed points besides the universal ceer degree and the universal  $\Sigma_2^0$ -degree?

For any  $\Delta_2^0$ -degree  $\mathbf{d}$ , the complete  $\mathbf{d}$ -ceer is a fixed-point. Are there any others? Is the universal ceer least among the fixed points?

Very little independent investigation here.

Many natural examples of things that correspond to ERs on  $2^\omega$  restricted to CE:  $=^{ce} \equiv \text{Id}^+ \in \Pi_2^0$ ,  $E_{set}^{ce} \equiv \text{Id}^{++} \in \Pi_4^0$ ,  $E_3^{ce} \in \Pi_4^0$

### Definition

For any  $E$ , let  $iE^+j$  if and only if  $[W_i]_E = [W_j]_E$ .

### Theorem

There is NO universal  $\Pi_n^0$ -equivalence relation.

In fact, for every  $\Pi_n^0$ -equivalence relation  $X$ , there is some  $\Delta_n^0$ -equivalence relation which is not below  $X$ .

This is a constant foot-gun. The temptation to say that  $=^{ce}$  is  $\Pi_2^0$ -universal is overpowering at times. Resist.

## Theorem

If  $X$  is a  $\Pi_2^0$ -equivalence relation, then there is some  $Y \in \Delta_2^0$  so that  $i X j$  iff  $Y^{[i]} = Y^{[j]}$ .

Now, the Ershov-Hierarchy essentially answers why there can't be a universal one. Consider the sequence:

$=^{ce}$  formed by letting  $Y$  be a universal c.e. set.

Next  $=^{d-ce}$  formed by letting  $Y$  be a universal d-c.e. set.

$\vdots$

$=^{\alpha-ce}$  formed by letting  $Y$  be a universal  $\alpha$ -c.e. set.

$\vdots$

By looking at where  $Y$  sits in the Ershov hierarchy, it's clear that these are co-final among  $\Delta_2^0$ -equivalence relations.

# Aside on $=\Sigma_n^0$ and $\dagger$

Relativizing at higher levels, that same hierarchy looks like:  
 $=\Sigma_3^0 < =\Sigma_{2n-1}^0 < \dots$

## Theorem

$$\text{Id}^{\dagger n} \equiv =\Sigma_{2n-1}^0.$$

## Corollary

Every  $\Sigma_{2n-1}^0$  or  $\Pi_{2n-1}^0$  equivalence relation reduces to  $\text{Id}^{\dagger n}$ .

## Proof.

If  $X$  is  $\Sigma_{2n-1}^0$ , we provide a reduction of  $X$  to  $=\Sigma_{2n-1}^0$ . Send  $n$  to  $[n]_X$ .

If  $X$  is  $\Pi_{2n-1}^0$ , send  $n$  to  $\omega \setminus [n]_X$ . □

## Question

Is  $\pi^{\dagger} \equiv =\Sigma_2^0$ ?



$\dagger$  doesn't preserve these difference hierarchies:

## Question

For any  $\Pi_n^0$ -equivalence relation  $X$ ,  $X^\dagger \leq_{\Sigma_{n+1}^0}$ .

## Proof.

Send  $i$  to  $[W_i]_X$ . □

## Question

We can ask about what the high  $\Pi_n^0$ -equivalence relations are. This has been looked at for the ceers with some surprising answers, but not even at  $\Pi_1^0$ .

Is  $=^{ce}$  the least  $\Pi_2^0$ -equivalence relation  $X$  so that  $X^\dagger \equiv_{\Sigma_3^0}$ ?  
Do they all have that jump?

# So why is there a $\Pi_1^0$ -universal?

## Theorem

For every  $\Pi_1^0$  relation (*not assumed transitive*)  $E$ , there is a  $\Delta_1^0$  set  $X$  and a partial computable function  $f$  so that if  $E$  is an equivalence relation, then  $i E j$  iff  $X^{[f(i)]} = X^{[f(j)]}$ .

## Proof.

At every  $s$ , we determine  $X(\langle n, m \rangle)$  for  $n, m \leq s$ . Let  $t_0 = 0$  and let  $t_{n+1}$  be the first stage  $> t_n$  where  $E$  looks transitive on  $[0, n+1]$ . If  $E$  is transitive, then this is an infinite sequence of stages, and  $f : n \mapsto t_n$  will be our reduction. When  $s$  is not a  $t_n$ -stage for some  $n$ , we do nothing much in coding  $X$  – make no differences. Put 0 on all new inputs.

Otherwise, code the highest-priority split – use transitivity to make all the coding columns look okay.  $\square$

We could do this for  $\Pi_2^0$ -relations, but the reduction function  $f$  would also be  $\Delta_2^0$ , so we wouldn't get computable reduction.

Here lie some natural ERs on c.e. sets:

$$E_0^{ce} \equiv E_1^{ce} \equiv E_2^{ce} \equiv \text{the } \Sigma_3^0\text{-universal degree}$$

### Definition

$iE_0^{ce} j$  iff  $W_i =^* W_j$

$iE_1^{ce} j$  iff for all but finitely many  $n$ ,  $W_i^{[n]} = W_j^{[n]}$

$iE_2^{ce} j$  iff  $\sum_{n \in A \Delta B} \frac{1}{n} < \infty$

The pattern seems to be that almost any “natural”  $\Sigma_n^0$ -equivalence relation will collapse to being universal. Obviously, this doesn’t happen at  $\Pi$ -levels.

Some classes within  $\Sigma_3^0$ -ERs, including the following two attempts to “effectivize” the class of countable borel equivalence relations (cbers).

# Countable Borel equivalence relations?

Definition (Coskey, Hamkins, R. Miller (2012))

- The action of a computable group  $G$  acting on  $\mathbf{CE}$  is **computable in indices** if there is computable  $\alpha$  so that

$$W_{\alpha(g,e)} = g \cdot W_e.$$

The induced orbit equivalence relation is denoted  $E_G^{ce}$ .

- $E^{ce}$  is **enumerable in indices** if there is computable  $\alpha$  so that, for all  $i \in \omega$ ,

$$e E^{ce} i \Leftrightarrow (\exists n)(W_{\alpha(e,n)} = W_i).$$

The first here was a natural attempt to use the Feldman-Moore theorem to bring the idea of cbers to ERs on  $\mathbf{CE}$ . The second attempt is similar, but using the Luzin-Novikov theorem.

## Theorem

If  $G$  is a computable group acting on  $\mathbf{CE}$  computably in indices, then either  $E_G^{ce} \equiv E_0^{ce}$  or  $E_G^{ce} \equiv_{=ce}$

First, we showed that any group acting on  $\mathbf{CE}$  computably in indices is actually acting via a permutation on  $\omega$ . Still, there are several computable subgroups of  $S_\infty$  to consider.

The prototypical examples to consider come down to the following cases:

- Let  $G$  be all finite permutations of  $\omega$ .
- Let  $\mathbb{Z}$  act on  $\omega$  by shifting.
- Let  $G$  be generated by  $(0, 1)(2, 3, 4)(5, 6, 7, 8) \cdots$ .

Having shown these were all  $\Sigma_3^0$ -complete, we realized that we had enough tricks to prove the same for any infinite  $G \subseteq S_\infty$ .

## Theorem

There are infinite chains and antichains of ERs which are enumerable in indices between  $=^{ce}$  and  $E_0^{ce}$ .

## Simple construction for chains.

For  $X \subsetneq \omega$ , let  $F(X)$  be the least element in  $X^c$ .

Let  $iR_nj$  if and only if  $W_i = W_j$  or  $0 \in W_i \cap W_j$  and  $F(W_i) \equiv F(W_j) \pmod n$ .

Note that  $=^{ce}$  reduces to  $R_n$  by sending  $W_i$  to  $W_i + 1$ . Among c.e. sets which contain 0, there are  $n + 1$  classes depending  $F(W_i) \pmod n$  OR  $F(W_i) = \infty$ . The last one is  $\Pi_2^0$ -complete, while the others are  $\Sigma_2^0$ -complete. By counting the number of properly  $\Sigma_2^0$ -classes, you can show  $R_{n+1} \not\leq R_n$ . □

Our examples are all  $\Delta_3^0$ . Can there be a properly  $\Sigma_3^0$ , but not universal, ER which is enumerable in indices?

Also, there is a  $\Delta_2^0$  enumerable in indices ER:  $E_{min}$ , and a  $\Pi_2^0$  which is below  $=^{ce}$ :  $E_{max}$ .

Can there be a  $\Sigma_2^0$  one which is not  $\Delta_2^0$ . More generally, can there be any  $\Sigma_2^0$  quotient of  $=^{ce}$  which is not  $\Delta_2^0$ ?

# Uniform enumeration in indices

Can the Lusin-Novikov direction be salvaged by demanding more uniformity from the enumerations?

## Definition

$E^{ce}$  is *uniformly* enumerable in indices if there is a computable  $\alpha$  so that for all  $i \in \omega$ ,

$$e E^{ce} i \Leftrightarrow (\exists n)(W_{\alpha(e,n)} = W_i).$$

and whenever  $W_e = W_i$ ,  $W_{\alpha(e,n)} = W_{\alpha(i,n)}$ .

Note that you expect this if the operation  $W_i \mapsto W_{\alpha(i,n)}$  is really an operation on sets (i.e., is independent of the enumeration).

## Observation

$E^{ce}$  is uniformly enumerable in indices if and only if it is the orbit equivalence of a computable action of a monoid  $M$  on **CE**.



for your attention, comments and contributions!