

A Jump in the Weihrauch degrees

Steffen Lemp

University of Wisconsin-Madison

March 19, 2024

(joint work with Andrews, Marcone, J. Miller and Valenti)

A *mathematical problem* can be viewed as a statement of the form

$$\forall X (\varphi(X) \rightarrow \exists Y \psi(X, Y)),$$

where φ and ψ are formulas in the (two-sorted) language $\mathcal{L} = \{+, \cdot, <, 0, 1, \in\}$ using only number quantifiers.

Here, X is called an *instance*, and Y a *solution* of the problem.

Two standard examples are:

- *Weak König's Lemma*: X is an infinite binary tree by $\varphi(X)$, and Y is an infinite path through X by $\psi(X, Y)$;
- *Ramsey's Theorem for Pairs and 2 Colors*: X is a 2-coloring of unordered pairs of numbers by $\varphi(X)$, and Y is an infinite homogeneous set by $\psi(X, Y)$.

This talk focuses on the Weihrauch degrees as a degree structure, rather than applications to measure the computational complexity of “ordinary” mathematical problems.

So will restrict ourselves to $X, Y \in \mathbb{N}^{\mathbb{N}}$ for the rest of the talk.

Consider problems as partial multi-valued functions $f : \subseteq \mathbb{N}^{\mathbb{N}} \rightrightarrows \mathbb{N}^{\mathbb{N}}$, mapping problems x satisfying $\varphi(x)$ to the set of all solutions y satisfying $\psi(x, y)$.

Definition (Weihrauch reducibility)

Let $f, g : \subseteq \mathbb{N}^{\mathbb{N}} \rightrightarrows \mathbb{N}^{\mathbb{N}}$ be partial multi-valued functions. Then f is *Weihrauch reducible* to g (written $f \leq_W g$) if the problem of computing f can be computably and uniformly solved by using for each instance of f a single computation of g in the following sense:

There are Turing functionals Φ and Ψ such that

- 1 $\Phi : \subseteq \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}^{\mathbb{N}}$ computes from every $p \in \text{dom}(f)$ some $q \in \text{dom}(g)$; and
- 2 $\Psi : \subseteq \mathbb{N}^{\mathbb{N}} \times \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}^{\mathbb{N}}$ computes from p and every $r \in g(q)$ some element of $f(p)$.

Basic Facts:

Let $\mathcal{W} = (\mathcal{PF} / \equiv_{\mathcal{W}}, \leq)$ (for the set \mathcal{PF} of partial multi-valued functions $f : \subseteq \mathbb{N}^{\mathbb{N}} \rightrightarrows \mathbb{N}^{\mathbb{N}}$), under the partial order induced by $\leq_{\mathcal{W}}$.

\mathcal{W} is a partial order with least element $\mathbf{0} = \{\emptyset\}$.

Under AC, \mathcal{W} has no greatest (or even maximal) element.

\mathcal{W} has size $2^{\mathfrak{c}} = 2^{2^{\aleph_0}}$.

In fact, every Weihrauch degree $\neq \mathbf{0}$ has size $2^{\mathfrak{c}}$.

Every nontrivial lower cone in \mathcal{W} has size $2^{\mathfrak{c}}$.

Every nontrivial maximal antichain in \mathcal{W} must be uncountable.

There is a maximal antichain of size $2^{\mathfrak{c}}$, but nothing more is known.

Every well-ordered ascending chain in \mathcal{W} of countable cofinality has an upper bound.

For every $\kappa \leq \mathfrak{c}$ of uncountable cofinality, there is an ascending chain in \mathcal{W} of type κ without upper bound.

(This is open for $\mathfrak{c} < \kappa \leq 2^{\mathfrak{c}}$.)

In 2012, Brattka, Gherardi and Marcone introduced an operation they called jump on partial multi-valued functions.

Their *jump* f' of f takes as input a sequence of elements of $\mathbb{N}^{\mathbb{N}}$ converging to some input for f ; and the output is the output of f .

For many f that they were interested then, they proved $f <_{\mathbb{W}} f'$, but especially at higher levels of the Weihrauch lattice, there are quite a few examples of interesting g such that $g' \equiv_{\mathbb{W}} g$:

So the jump sometimes does not “jump” (i.e., strictly increase)!

Moreover, the jump is not degree-theoretic:

There exist f and g such that $f \equiv_{\mathbb{W}} g$ while $f' \not\equiv_{\mathbb{W}} g'$.

So the jump is not *weakly monotone*: There exist f and g such that $f \leq_{\mathbb{W}} g$ while $f' \not\leq_{\mathbb{W}} g'$.

(However, the jump is weakly monotone with respect to *strong* Weihrauch reducibility.)

So Marcone and Valenti started looking for a “better-behaved” jump operation.

For a partial order $\mathcal{P} = (P, \leq_P)$, we define a *jump operation* on \mathcal{P} to be a function $j : P \rightarrow P$ that is

- 1 *strictly increasing*, i.e., $p <_P j(p)$ for every $p \in P$; and
- 2 *weakly monotone*, i.e., if $p \leq_P q$ then $j(p) \leq_P j(q)$ for all $p, q \in P$.

Using the Axiom of Choice, one can show that every upper semilattice without maximum, and in particular the Weihrauch lattice, has an “abstract” jump operation.

But we wanted a “natural” jump operation on the Weihrauch lattice that satisfies the above conditions.

Note: In preliminary work with Raghavan, I have a sketch of an argument that ZF alone does not suffice to ensure a jump operation on every upper semilattice without maximum.

We first need the following

Definition (totalization)

The *totalization* (or *total continuation*) of $f : \subseteq \mathbb{N}^{\mathbb{N}} \rightrightarrows \mathbb{N}^{\mathbb{N}}$ is the total (multi-valued) function $\mathbb{T}f : \mathbb{N}^{\mathbb{N}} \rightrightarrows \mathbb{N}^{\mathbb{N}}$ defined by

$$\mathbb{T}f(p) = \begin{cases} f(p) & \text{if } p \in \text{dom}(f); \\ \mathbb{N}^{\mathbb{N}} & \text{otherwise.} \end{cases}$$

The totalization has been before used; e.g., the Weihrauch degree of $\text{TC}_{\mathbb{N}^{\mathbb{N}}}$ is relevant for functions related to the system ATR_0 .

The totalization is not a degree-theoretic operation and, even for non-total f , it is possible that $\mathbb{T}f \equiv_{\text{W}} f$.

The totalization comes close to giving a natural jump operation, but it requires a small tweak:

Definition (totalizing jump)

The Weihrauch degree of the *totalizing jump*, or *tot-jump*, of f is the maximum of the degrees of $\top g$ for all $g \equiv_W f$.

(We denote it by $\text{tJ}(f)$.)

This maximum exists since it can also be defined as follows:

For $x = (e, i) \frown p \in \mathbb{N}^{\mathbb{N}}$, define

$$\text{tJ}(f)(x) = \begin{cases} \{ \Phi_i(p, q) \mid q \in f\Phi_e(p) \} & \text{if } \Phi_e(p) \in \text{dom}(f) \text{ and} \\ & (\forall q \in f\Phi_e(p))(\Phi_i(p, q) \downarrow) \\ \mathbb{N}^{\mathbb{N}} & \text{otherwise} \end{cases}$$

This is indeed a jump operation:

Theorem

- 1 For every f , $f <_W \text{tJ}(f)$.
- 2 For every f and g , if $f \leq_W g$ then $\text{tJ}(f) \leq_W \text{tJ}(g)$.

Proof Sketch: $f \leq_W \text{tJ}(f)$ and weak monotonicity are easy.

To show $\text{tJ}(f) \not\leq_W f$, define first $d : \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}^{\mathbb{N}}$ by

$$d(p)(0) := p(0) + 1, \text{ and } d(p)(n) := p(n) \text{ for } n > 0.$$

Then define a new function $f_d \equiv_W \text{tJ}(f)$ (this requires a proof!) by

$$f_d(x) := \begin{cases} \{d\Phi_i(x, q) \mid q \in f\Phi_e(x)\} & \text{if } \Phi_e(x) \in \text{dom}(f) \text{ and} \\ & (\forall q \in f\Phi_e(x))(\Phi_i(x, q) \downarrow), \\ \mathbb{N}^{\mathbb{N}} & \text{otherwise.} \end{cases}$$

Suppose $f_d \leq_W f$ via Φ_e and Φ_i . For $p \in \mathbb{N}^{\mathbb{N}}$, let $y = (e, i) \hat{\ } p$.

So $y \in \text{dom}(f_d)$; $\Phi_e(y) \in \text{dom}(f)$; and $(\forall q \in f\Phi_e(y))(\Phi_i(y, q) \downarrow)$.

Now set $X = \{\Phi_i(y, t) \mid t \in f\Phi_e(y)\} \neq \emptyset$, so there is $q \in f\Phi_e(y)$ with $\Phi_i(y, q) \notin d(X) = \{d\Phi_i(y, t) \mid t \in f\Phi_e(y)\} = f_d(y)$. \square

Another characterization of the totalizing jump is as follows:

Definition

Let $W_{\Pi_2^0 \rightarrow \Pi_1^0} \subseteq \mathbb{N}^{\mathbb{N}} \rightrightarrows \mathbb{N}^{\mathbb{N}}$ be the problem defined as

$$W_{\Pi_2^0 \rightarrow \Pi_1^0}(p) := \{q \in \mathbb{N}^{\mathbb{N}} \mid \forall i (q(i+1) > q(i) \text{ and } p(q(i)) = 0)\}.$$

So the domain of $W_{\Pi_2^0 \rightarrow \Pi_1^0}$ is the set of all functions p with infinitely many zeroes, and the outputs are increasing listings of positions of infinitely many of these zeroes.

(So $W_{\Pi_2^0 \rightarrow \Pi_1^0}$ converts a Π_2^0 -question into a Π_1^0 -question.)

Theorem

For every f , we have $\text{tJ}(f) \equiv_W \text{T}(W_{\Pi_2^0 \rightarrow \Pi_1^0} * f * W_{\Pi_2^0 \rightarrow \Pi_1^0})$.
(Here, $g * h$ is the sequential product of h and then g .)

But $W_{\Pi_2^0 \rightarrow \Pi_1^0}$ cannot be dropped on either side of f !

We have far more than for “usual jumps” on degree structures:

Theorem

*For every f, g , $f \leq_W g$ if and only if $tJ(f) \leq_W tJ(g)$.
Thus tJ is an (injective!) endomorphism on the Weihrauch degrees.*

Here are some specific examples of totalizing jumps:

Theorem

- $tJ(\emptyset) = \text{id}$;
- $tJ(\text{id})(p) = \{ q \mid \exists m \forall n > m p(n) = 1 \iff \exists n q(n) > 0 \}$;
- $tJ(C_{\mathbb{N}^{\mathbb{N}}}) = TC_{\mathbb{N}^{\mathbb{N}}}$.

*Moreover, $\text{LPO} <_W tJ(\text{id})$ and $\widehat{tJ(\text{id})} \equiv_W \text{lim}$.
(Here $\text{LPO} : \mathbb{N}^{\mathbb{N}} \rightarrow \{0, 1\}$ is defined by $\text{LPO}(p) = 0$ iff $p \neq 0^{\mathbb{N}}$;
and \hat{f} is the countable parallelization of f .)*

We also have explicit descriptions of $tJ^n(\text{id})$ for every n .

Main Open Question

Characterize the range of the totalizing jump on the Weihrauch degrees.

We are very far from a satisfactory answer.

In fact, while our definition of the totalizing jump is Δ_2^1 , we would like to know if we can lower that complexity to Π_1^1 , say, given that we can prove it cannot be Σ_1^1 .

But we do have some interesting properties of the totalizing jump:

Theorem

For every f , $tJ(f)$ is total and join-irreducible but can be meet-reducible.

However, not all total join-irreducible Weihrauch degrees are in the range of tJ , e.g., (the degree of) $TC_{\mathbb{N}}$ is not in the range of tJ .

Here are some more properties of the totalizing jump:

Theorem

For every $g \neq \emptyset$, there exists $f <_W g$ such that $tJ(f) \not\leq_W g$.

Theorem

For every f , there exists h s.t. $f <_W h <_W tJ(f)$.

Theorem

If $LPO \times g \leq_W tJ(f)$, then $g \leq_W f$.

In fact (with Pauly): If $DIS \times g \leq_W tJ(f)$, then $g \leq_W f$.

(Here, $DIS : \mathbb{N}^{\mathbb{N}} \rightrightarrows \mathbb{N}^{\mathbb{N}}$, $p \mapsto \{q \mid U(p) \neq q\}$ for a fixed universal partial computable functional $U : \subseteq \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}^{\mathbb{N}}$.)

Corollary

If $DIS \times g \leq_W g$, then g is not in the range of tJ .

Thank you!